



TRUST  SOFT
Mathematical Guarantees Eliminate Software Risk



TrustInSoft
Updates about the Frama-C
software publisher company



Benjamin Monate
co-founder & CTO

Company Facts

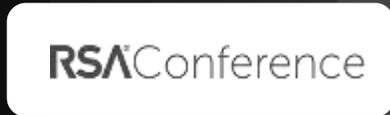
- French startup created in 2013 as a Spin-off of CEA



Only company selected in the Ockham Criteria from the SATE V exhibit



Chosen by the Linux Foundation to develop tools for security of Core Internet Infrastructure



Nominated as one the 10 most innovative companies in cybersecurity – RSA '15 Conference



Most Innovative SME, Special Jury prize 2016 at the Forum International de la Cybersécurité 2016 sponsored by Airbus Defence&Space

TrustInSoft Unique Value Proposal

provide guarantees
on software used in
sensitive systems

Current Customers

TrustInSoft works with the most demanding developers of sensitive software.

Since 2013



Aeronautics

DO-178C - ED-12C



Nuclear Reactors

IEC-60880 IEC-62138



Defense

Since 2014



Rail

EN-50128



Space



Telecom

Since 2015



Automotive

ISO 2626-2



Smart Factories



Cyber

CWE

Customer names are under strict NDAs - 50% in the US

ARM mbed can claim they have the first ever TLS/SSL stack without buffer overflows.

ARMmbed™

Using TrustInSoft Analyzer we have generated a report which tells how to compile, configure and deploy mbed TLS in a given perimeter in order to be immune from all attacks caused by CWE 119 to 127, 369, 415, 416, 457, 476, 562, 690.

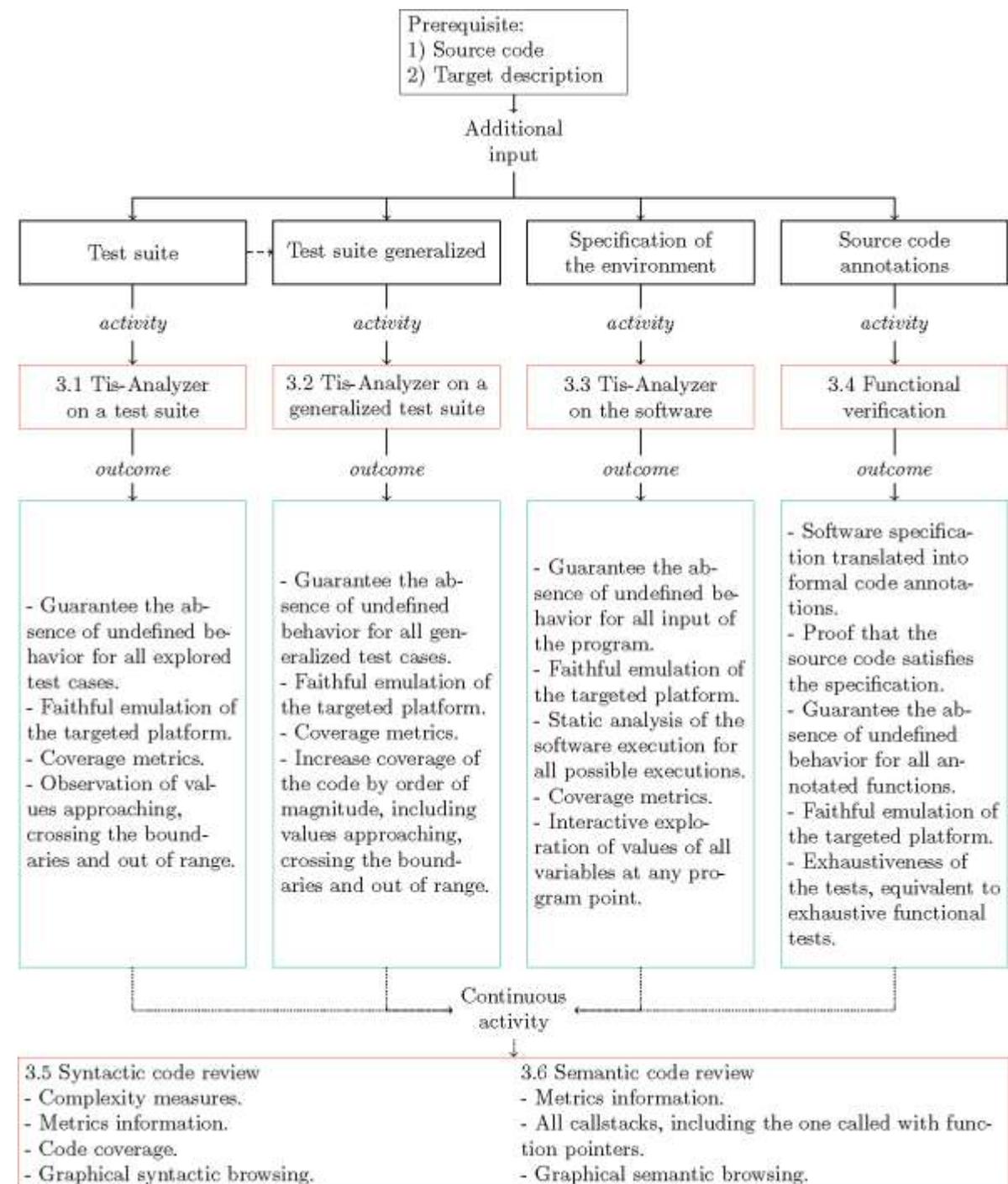


This stack has a configuration proven to be without an Heartbleed-like flaw.

You can download such a report here: <http://trust-in-soft.com/polarssl-verification-kit>

Progressive methodologies

- From enhanced testing to full functional verification
- Adapt level of verification to each customer
- Incremental guarantees



Standard compliance

Documents for common software verification standards

- ISO 26262,
- EN50128,
- DO178B/C,
- BV-SW-100,
- SEI CERT C
- Preparation phase: for each verification activity, state the support level customer can expect from TIS-Analyzer
- Production phase: propose specific tool adaptations to match the customer process

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Usage of Tis-Analyzer platform for
BV-SW-100

Version:

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Draft – Usage of Tis-Analyzer platform for
SEI CERT C Coding Standard

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Usage of Tis-Analyzer platform for
ISO-26262

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Usage of Tis-Analyzer platform for EN-50128

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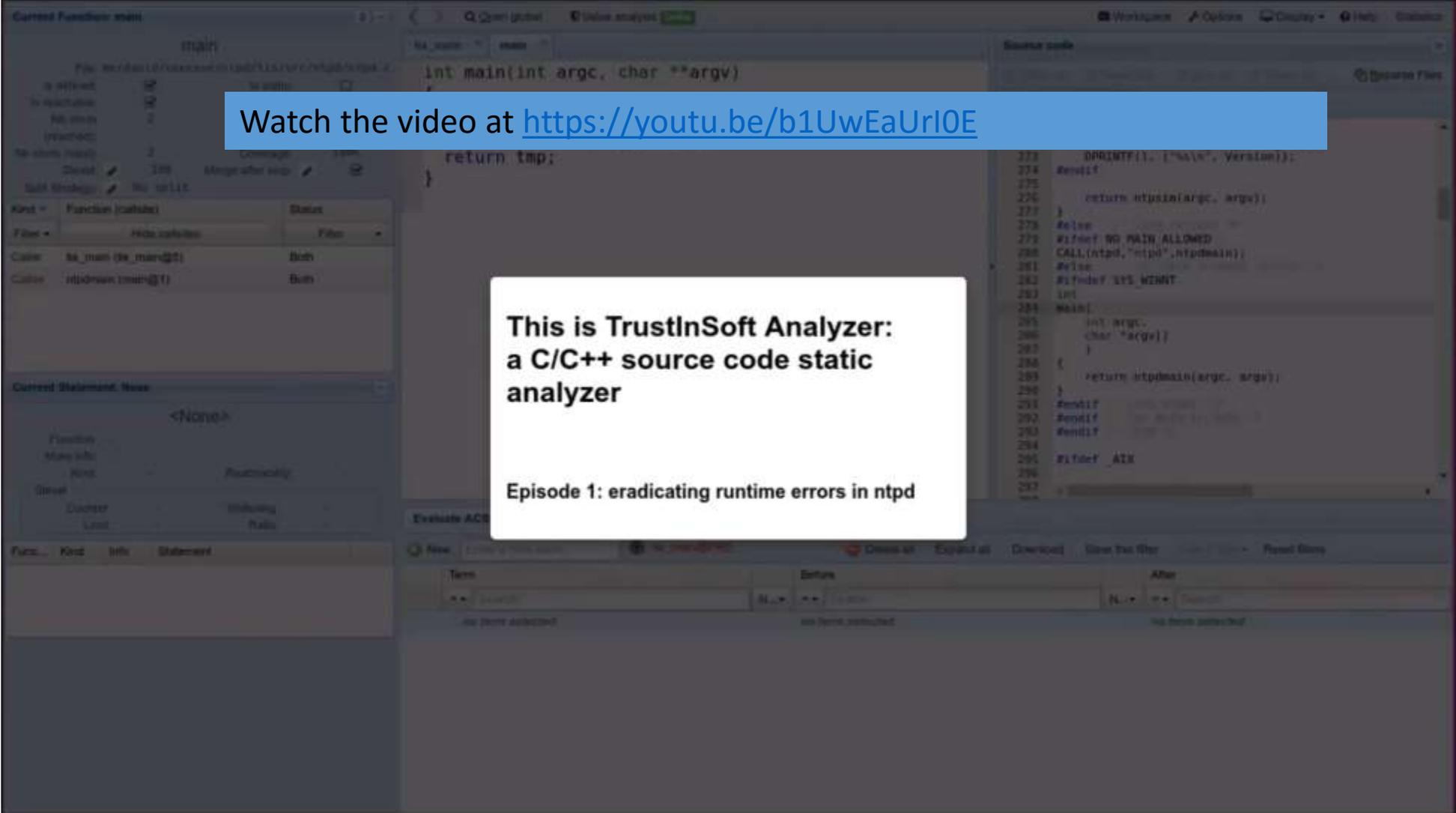
Usage of Tis-Analyzer platform for CWE

Version: 1.22

Training sessions: 1 to 5 days for engineers

- Introduction to formal method for software analysis
- Why undefined behaviors matter in C11
- TrustInSoft Analyzer for Safety Standard Compliance
- Testing Software with a perfect source semantics: TIS-Interpreter
- Eradicating undefined behaviors in existing C applications
- Developing secured applications in C with TIS-Analyzer

A glance at TrustInSoft Analyzer User Interface



Watch the video at <https://youtu.be/b1UwEaUrI0E>

**This is TrustInSoft Analyzer:
a C/C++ source code static
analyzer**

Episode 1: eradicating runtime errors in ntpd

The screenshot displays the TrustInSoft Analyzer interface. The central pane shows the source code of a C program, including a `main` function and a call to `ntpmain`. The left sidebar contains a project tree and a list of functions. The bottom section shows a table for function calls, with columns for 'Term', 'Before', and 'After'. The interface is dark-themed and includes various toolbars and status bars.

Innovation at TrustInSoft

- Global Dataflow Analysis

Seeding a random number generator

- Side Channel Attacks Analysis

Constant time or memory access

- Strict aliasing analysis

Typed memory model

Innovation at TrustInSoft

- Precision improvement for derived analysis

Value analysis result stored as a graph

- Analysis of large Open Source Software

Communication stack, compression libraries, image libraries

- Tis-Interpreter

Combination with state-of-the-art fuzzers

Open source at <https://github.com/TrustInSoft/tis-interpreter>

Innovation at TrustInSoft

TrustInSoft Analyzer for C++

Derived from CEA STANCE prototype

All C++11 language including lambdas

C++11 STL support

C++14 support

Support for ACSL++

Innovation at TrustInSoft: research projects

ANR ANASTASEC – Cyber-security in avionics

FUI INGOPCS – Secured Industrial Communication Stack

DGE S3P – Smart Safe and Secure Platform for IoT

FUI SecureOCaml – Securing OCaml programs

RAPID AUROCHS – Code Analysis of Cryptographic Stack

ANR VOCAL – Certified OCaml libraries

Any questions?

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Mathematical Guarantees Eliminate Software Risk

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BN_consttime_swap from OpenSSL: constant time

```
void BN_consttime_swap(BN_ULONG condition, BIGNUM *a,
BIGNUM *b, int nwords) {
    BN_ULONG t;
    int i;

    bn_wcheck_size(a, nwords);
    bn_wcheck_size(b, nwords);

    assert(a != b);
    assert((condition & (condition - 1)) == 0);
    assert(sizeof(BN_ULONG) >= sizeof(int));

    condition = ((condition - 1) >> (BN_BITS2 - 1)) - 1;

    t = (a->top ^ b->top) & condition;
    a->top ^= t;
    b->top ^= t;

#define BN_CONSTTIME_SWAP(ind) \
    do { \
        t = (a->d[ind] ^ b->d[ind]) & condition; \
        a->d[ind] ^= t; \
        b->d[ind] ^= t; \
    } while (0)

    switch (nwords) {
    default:
        for (i = 10; i < nwords; i++)
            BN_CONSTTIME_SWAP(i);
        /* Fallthrough */
    case 10:
        BN_CONSTTIME_SWAP(9); /* Fallthrough */
    case 9:
        BN_CONSTTIME_SWAP(8); /* Fallthrough */
    case 8:
        BN_CONSTTIME_SWAP(7); /* Fallthrough */
    case 7:
        BN_CONSTTIME_SWAP(6); /* Fallthrough */
    case 6:
        BN_CONSTTIME_SWAP(5); /* Fallthrough */
    case 5:
        BN_CONSTTIME_SWAP(4); /* Fallthrough */
    case 4:
        BN_CONSTTIME_SWAP(3); /* Fallthrough */
    case 3:
        BN_CONSTTIME_SWAP(2); /* Fallthrough */
    case 2:
        BN_CONSTTIME_SWAP(1); /* Fallthrough */
    case 1:
        BN_CONSTTIME_SWAP(0);
    }
#undef BN_CONSTTIME_SWAP
}
```

```
switch (nwords) {
    default:
        for (i = 10; i < nwords; i++)
            BN_CONSTTIME_SWAP(i);
        /* Fallthrough */
    case 10:
        BN_CONSTTIME_SWAP(9); /* Fallthrough */
    case 9:
        BN_CONSTTIME_SWAP(8); /* Fallthrough */
    case 8:
        BN_CONSTTIME_SWAP(7); /* Fallthrough */
    case 7:
        BN_CONSTTIME_SWAP(6); /* Fallthrough */
    case 6:
        BN_CONSTTIME_SWAP(5); /* Fallthrough */
    case 5:
        BN_CONSTTIME_SWAP(4); /* Fallthrough */
    case 4:
        BN_CONSTTIME_SWAP(3); /* Fallthrough */
    case 3:
        BN_CONSTTIME_SWAP(2); /* Fallthrough */
    case 2:
        BN_CONSTTIME_SWAP(1); /* Fallthrough */
    case 1:
        BN_CONSTTIME_SWAP(0);
    }
#undef BN_CONSTTIME_SWAP
}
```

We automatically confirm there are no problems in the BN_consttime_swap function

s2n_verify_cbc from S2N: constant time?

```
int s2n_verify_cbc(struct s2n_connection *conn,
                  struct s2n_hmac_state *hmac,
                  struct s2n_blob *decrypted)
{
    // . . . .

    int cutoff = check - padding_length;
    for (int i = 0, j = decrypted->size - 1 - check;
         i < check && j < decrypted->size; i++, j++) {
        uint8_t mask = ~(0xff << ((i >= cutoff) * 8));
        mismatches |= (decrypted->data[j] ^ padding_length) & mask;
    }
    if (mismatches) {
        S2N_ERROR(S2N_ERR_CBC_VERIFY);
    }
    return 0;
}
```

This function has been written to compute if there is a padding problem in the variable `mismatches` and to return this value.

When not in debug mode the `S2N_ERROR` function returns simply the error code. A possible solution
`return (!!mismatches) * S2N_ERR_CBC_VERIFY;`

Error detected

Detected undefined behaviors

- Division by zero
- Memory Accesses
 - Memory access
 - Index out of bound
 - Memory problem
 - Overflow in array accesses
- Valid string
- Invalid shift
- Pointer comparison
- Differing blocks
- Overflow
- Float to integer
- Not separated
- Overlap
- Dangling and Uninitialized pointers
 - Initialization
 - Dangling
- Is nan or infinite
- Function type matches

Other possible analyses

- Exhaustive detection of all undefined behaviors of the program
- Functional dependency analyses
- Control and data flow analyses
- Shared variable detection
- Race condition detection
- Memory leak
- Proof of functional properties
- Command line and graphical interface

Dedicated all-included support

- All technical questions answered
- Tool Expertise and Customization
- Review of your integration processes and normative requirements