



Software Analyzers

## E-ACSL Version 1.16

Implementation in Frama-C plug-in E-ACSL  
version 23.0 rc1







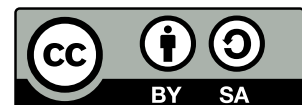
# E-ACSL

## Executable ANSI/ISO C Specification Language

Version 1.16 – Frama-C plug-in E-ACSL version 23.0 rc1

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# Foreword

This document describes version 1.16 of the E-ACSL specification language. It is based on the ACSL specification language [1]. Features of both languages may still evolve in the future, even if we do our best to preserve backward compatibility.

## Acknowledgements

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# Chapter 1

## Introduction

This document is a reference manual for the E-ACSL implementation provided by the E-ACSL plug-in [10] (version 23.0 rc1) of the FRAMA-C framework [5]. E-ACSL is an acronym for “Executable ANSI/ISO C Specification Language”. It is an “executable” subset of *stable* ACSL [1] implemented [2] in the FRAMA-C platform [5]. “Stable” means that no experimental ACSL feature is supported by E-ACSL. Contrary to ACSL, each E-ACSL specification is executable: it may be evaluated at runtime.

In this document, we assume that the reader has a good knowledge of both ACSL [1] and the ANSI C programming language [7, 8].

### 1.1 Organization of this document

---

This document is organized in the very same way that the reference manual of ACSL [1].

Instead of being a fully new reference manual, this document points out the differences between E-ACSL and ACSL. Each E-ACSL construct which is not pointed out must be considered to have the very same semantics than its ACSL counterpart. For clarity, each relevant grammar rules are given in BNF form in separate figures like the ACSL reference manual does. In these rules, constructs with semantic changes are displayed in blue.

Not all of the features mentioned in this document are currently implemented in the FRAMA-C’s E-ACSL plug-in. Those who aren’t yet are signaled as in the following line:

This feature is not currently supported by FRAMA-C’s E-ACSL plug-in.<sup>1</sup>

As a summary, Figure 1.1 synthetizes main features that are not currently implemented into the FRAMA-C’s E-ACSL plug-in.

### 1.2 Generalities about Annotations

---

*No difference with ACSL.*

### 1.3 Notations for grammars

---

*No difference with ACSL.*

---

<sup>1</sup>Additional remarks on the feature may appear as footnote.

typing	mathematical reals
terms	<code>\true</code> and <code>\false</code> let binding t-sets
predicates	exclusive or operator let bindings quantifications over non-integer types <code>\specified</code>
annotations	behavior-specific annotations loop assigns global annotations
behavior clauses	assigns allocates decreases abrupt termination

Figure 1.1: Summary of not-yet-implemented features.

## Specification language

## 2.1 Lexical rules

---

*No difference with ACSL.*

## 2.2 Logic expressions

---

*No difference with ACSL, but guarded quantification..*

More precisely, grammars of terms and binders presented respectively Figures 2.1 and 2.3 are the same than the one of ACSL, while Figure 2.2 presents grammar of predicates. The only difference between E-ACSL and ACSL predicates are quantifications.

### Quantification

E-ACSL quantification must be computable. They are limited to two limited forms.

**Guarded integer quantification** Guarded universal quantification is denoted by

```
\forall x_1, \dots, x_n;
  a_1 <= x_1 <= b_1 ... && a_n <= x_n <= b_n
  ==> p
```

and guarded existential quantification by

```
\exists x_1, \dots, x_n;
  a_1 <= x_1 <= b_1 ... && a_n <= x_n <= b_n
  && p
```

Each variable must be guarded exactly once and the guard of  $x_i$  must appear before the guard of  $x_j$  if  $i < j$  (*i.e.* order of guards must follow order of binders).

Following the definition, each quantified variable belongs to a finite interval. Since finite interval is only computable in practice for integers, this form of quantifier is limited to `integer` and its subtype. Thus there is no guarded quantification over `float`, `real`, `C` pointers or logic types.

<i>literal</i>	<code>::=</code>	<code>\true</code>   <code>\false</code>   <code>integer</code>   <code>real</code>   <code>string</code>   <code>character</code>	boolean constants integer constants real constants string constants character constants
<i>bin-op</i>	<code>::=</code>	<code>+</code>   <code>-</code>   <code>*</code>   <code>/</code>   <code>%</code>   <code>&lt;&lt;</code>   <code>&gt;&gt;</code>   <code>==</code>   <code>!=</code>   <code>&lt;=</code>   <code>&gt;=</code>   <code>&gt;</code>   <code>&lt;</code>   <code>&amp;&amp;</code>   <code>  </code>   <code>^^</code>   <code>&amp;</code>   <code> </code>   <code>--&gt;</code>   <code>&lt;--&gt;</code>   <code>^</code>	boolean operations bitwise operations
<i>unary-op</i>	<code>::=</code>	<code>+</code>   <code>-</code>   <code>!</code>   <code>~</code>   <code>*</code>   <code>&amp;</code>	unary plus and minus boolean negation bitwise complementation pointer dereferencing address-of operator
<i>term</i>	<code>::=</code>	<code>literal</code>   <code>id</code>   <code>unary-op term</code>   <code>term bin-op term</code>   <code>term [ term ]</code>   <code>{ term \with [ term ] = term }</code>   <code>term . id</code>   <code>{ term \with . id = term }</code>   <code>term -&gt; id</code>   <code>( type-expr ) term</code>   <code>id ( term ( , term)* )</code>   <code>( term )</code>   <code>term ? term : term</code>   <code>\let id = term ; term</code>   <code>sizeof ( term )</code>   <code>sizeof ( C-type-name )</code>   <code>id : term</code>   <code>string : term</code>	literal constants variables  array access array functional modifier structure field access field functional modifier  cast function application parentheses ternary condition local binding  syntactic naming syntactic naming

Figure 2.1: Grammar of terms

<i>rel-op</i>	::=	<code>==   !=   &lt;=   &gt;=   &gt;   &lt;</code>	
<i>pred</i>	::=	<code>\true   \false</code>	
		<code>term (rel-op term)<sup>+</sup></code>	comparisons
		<code>id ( term ( , term)<sup>*</sup> )</code>	predicate application
		<code>( pred )</code>	parentheses
		<code>pred &amp;&amp; pred</code>	conjunction
		<code>pred    pred</code>	disjunction
		<code>pred ==&gt; pred</code>	implication
		<code>pred &lt;==&gt; pred</code>	equivalence
		<code>! pred</code>	negation
		<code>pred ^^ pred</code>	exclusive or
		<code>term ? pred : pred</code>	ternary condition
		<code>pred ? pred : pred</code>	
		<code>\let id = term ; pred</code>	local binding
		<code>\let id = pred ; pred</code>	
		<code>\forallall binders ;</code>	
		<code>integer-guards ==&gt; pred</code>	univ. integer quantification
		<code>\exists binders ;</code>	
		<code>integer-guards &amp;&amp; pred</code>	exist. integer quantification
		<code>\forallall binders ;</code>	
		<code>iterator-guard ==&gt; pred</code>	univ. iterator quantification
		<code>\exists binders ;</code>	
		<code>iterator-guard &amp;&amp; pred</code>	exist. iterator quantification
		<code>\forallall binders ; pred</code>	univ. quantification
		<code>\exists binders ; pred</code>	exist. quantification
		<code>id : pred</code>	syntactic naming
		<code>string : pred</code>	syntactic naming
<i>integer-guards</i>	::=	<code>interv (&amp;&amp; interv)<sup>*</sup></code>	
<i>interv</i>	::=	<code>(term integer-guard-op)<sup>+</sup></code>	
		<code>id</code>	
		<code>(integer-guard-op term)<sup>+</sup></code>	
<i>integer-guard-op</i>	::=	<code>&lt;=   &lt;</code>	
<i>iterator-guard</i>	::=	<code>id ( term , term )</code>	

Figure 2.2: Grammar of predicates

$binders$	$::=$	$binder$	$(, binder)^*$	
$binder$	$::=$	$type\text{-}expr$	$variable\text{-}ident$	
			$(, variable\text{-}ident)^*$	
$type\text{-}expr$	$::=$	$logic\text{-}type\text{-}expr$	$ $	$C\text{-}type\text{-}name$
$logic\text{-}type\text{-}expr$	$::=$	$built\text{-}in\text{-}logic\text{-}type$		
		$ $	$id$	type identifier
$built\text{-}in\text{-}logic\text{-}type$	$::=$	$boolean$	$ $	$integer$
			$ $	$real$
$variable\text{-}ident$	$::=$	$id$	$ $	$* variable\text{-}ident$
			$ $	$variable\text{-}ident$ $[]$
			$ $	$( variable\text{-}ident )$

Figure 2.3: Grammar of binders and type expressions

**Iterator quantification** In order to iterate over non-integer types, E-ACSL introduces a notion of *iterators* over types: standard ACSL unguarded quantifications are only allowed over a type which an iterator is attached to.

Iterators are introduced by a specific construct which attaches two sets — namely **nexts** and the **guards** — to a binary predicate over a type  $\tau$ . Both sets must have the same cardinal. This construct is described by the grammar of Figure 2.4. For a type  $\tau$ ,

$declaration$	$::=$	$//@$	$iterator$	$id$	$($	$wildcard\text{-}param$	$,$	$wildcard\text{-}param$	$)$	$:$
						$nexts$		$terms$		$;$
						$guards$		$predicates$		$;$
$wildcard\text{-}param$	$::=$	$parameter$								
			$ $			$-$				
$terms$	$::=$	$term$	$(,$	$term$	$)^*$					
$predicates$	$::=$	$predicate$	$(,$	$predicate$	$)^*$					

Figure 2.4: Grammar of iterator declarations

**nexts** is a set of terms which take an argument of type  $\tau$  and return a value of type  $\tau$  which computes the next element in this type, while **guards** is a set of predicates which take an argument of type  $\tau$  and are valid (resp. invalid) to continue (resp. stop) the iteration.

Furthermore, the guard of a quantification using an iterator must be the predicate given in the definition of the iterator. This abstract binary predicate takes two arguments of the same type. One of them must be unnamed by using a wildcard (character underscore '\_'). The unnamed argument must be binded to the quantifier, while the other corresponds to the term from which the iteration begins.

**Example 2.1** *The following example introduces binary trees and a predicate which is valid if and only if each value of a binary tree is even.*

```
struct btree {
  int val;
  struct btree *left, *right;
};
```



```

/*@ iterator access(_, struct btree *t):
  @ nexts t->left, t->right;
  @ guards \valid(t->left), \valid(t->right); */

/*@ predicate is_even(struct btree *t) =
  @ \forall struct btree *tt; access(tt, t) ==> tt->val % 2 == 0; */

```

**Unguarded quantification** They are only allowed over boolean and char.

### 2.2.1 Operators precedence

*No difference with ACSL.*

Figure 2.5 summarizes operator precedences.

class	associativity	operators
selection	left	[...] -> .
unary	right	! ~ + - * & (cast) sizeof
multiplicative	left	* / %
additive	left	+ -
shift	left	<< >>
comparison	left	< <= > >=
comparison	left	== !=
bitwise and	left	&
bitwise xor	left	^
bitwise or	left	
bitwise implies	left	-->
bitwise equiv	left	<-->
connective and	left	&&
connective xor	left	^^
connective or	left	
connective implies	right	==>
connective equiv	left	<==>
ternary connective	right	...?...:...>
binding	left	\forall \exists \let
naming	right	:

Figure 2.5: Operator precedence

### 2.2.2 Semantics

*No difference with ACSL, but undefinedness and same laziness than C.*

More precisely, while ACSL is a 2-valued logic with only total functions, E-ACSL is a 3-valued logic with partial functions since terms and predicates may be “undefined”.

In this logic, the semantics of a term denoting a C expression  $e$  is undefined if  $e$  leads to a runtime error. Consequently the semantics of any term  $t$  (resp. predicate  $p$ ) containing a C expression  $e$  leading to a runtime error is undefined if  $e$  has to be evaluated in order to evaluate  $t$  (resp.  $p$ ).

**Example 2.2** *The semantics of all the below predicates are undefined:*

- $1/0 == 1/0$
- $f(*p)$  for any logic function  $f$  and invalid pointer  $p$

Furthermore, C-like operators  $\&\&$ ,  $\|\|$ ,  $\wedge\wedge$  and  $\_ ? \_ : \_$  are lazy like in C: their right members are evaluated only if required. Thus the amount of undefinedness is limited. Consequently, predicate  $p ==> q$  is also lazy since it is equivalent to  $!p \|\| q$ . It is also the case for guarded quantifications since guards are conjunctions and for ternary condition since it is equivalent to a disjunction of implications.

**Example 2.3** *Below, the first, second and fourth predicates are invalid while the third one is valid:*

- $\backslash\text{false} \ \&\& \ 1/0 == 1/0$
- $\backslash\text{forall integer } x, -1 \leq x \leq 1 ==> 1/x > 0$
- $\backslash\text{forall integer } x, 0 \leq x \leq 0 ==> \backslash\text{false} ==> -1 \leq 1/x \leq 1$
- $\backslash\text{exists integer } x, 1 \leq x \leq 0 \ \&\& \ -1 \leq 1/x \leq 1$

*In particular, the second one is invalid since the quantification is in fact an enumeration over a finite number of elements, it amounts to  $1/-1 > 0 \ \&\& \ 1/0 > 0 \ \&\& \ 1/1 > 0$ . The first atomic proposition is invalid, so the rest of the conjunction (and in particular  $1/0$ ) is not evaluated. The fourth one is invalid since it is an existential quantification over an empty range.*

*A contrario the semantics of predicates below is undefined:*

- $1/0 == 1/0 \ \&\& \ \backslash\text{false}$
- $-1 \leq 1/0 \leq 1 ==> \backslash\text{true}$
- $\backslash\text{exists integer } x, -1 \leq x \leq 1 \ \&\& \ 1/x > 0$

Furthermore, casting a term denoting a C expression  $e$  to a smaller type  $\tau$  is undefined if  $e$  is not representable in  $\tau$ .

**Example 2.4** *Below, the first term is well-defined, while the second one is undefined.*

- $(\text{char})127$
- $(\text{char})128$

**Handling undefinedness in tools** It is the responsibility of each tool which interprets E-ACSL to ensure that an undefined term is never evaluated. For instance, they may exit with a proper error message or, if they generate C code, they may guard each generated undefined C expression in order to be sure that they are always safely used.

This behavior is consistent with both ACSL [1] and mainstream specification languages for runtime assertion checking like JML [9]. Consistency means that, if it exists and is defined, the E-ACSL predicate corresponding to a valid (resp. invalid) ACSL predicate is valid (resp. invalid). Thus it is possible to reuse tools interpreting ACSL like the FRAMA-C's value analysis plug-in [6] in order to interpret E-ACSL, and it is also possible to perform runtime assertion checking of E-ACSL predicates in the same way than JML predicates. Reader interested by the implications (especially issues) of such a choice may read articles of Patrice Chalin [3, 4].

### 2.2.3 Typing

*No difference with ACSL, but no user-defined types.*

It is not possible to define logic types introduced by the specification writer (see Section 2.6).

### 2.2.4 Integer arithmetic and machine integers

*No difference with ACSL.*

### 2.2.5 Real numbers and floating point numbers

*No difference with ACSL.*

*Exact real numbers and even floating point numbers are usually difficult to implement. Thus you would not wonder if most tools do not support them (or support them partially).*

**Real numbers beyond rationals are currently not supported by the E-ACSL plug-in. Only rationals (in  $\mathbb{Q}$ ) and floating point numbers are supported.**

### 2.2.6 C arrays and pointers

*No difference with ACSL.*

*Ensuring validity of memory accesses is usually difficult to implement, since it requires the implementation of a memory model. Thus you would not wonder if most tools do not support it (or support it partially).*

### 2.2.7 Structures, Unions and Arrays in logic

*No difference with ACSL.*

*Logic arrays without an explicit length are usually difficult to implement. Thus you would not wonder if most tools do not support them (or support them partially).*

**The `\length` function is currently not supported by the E-ACSL plug-in.**

**The comparison of unions and structures is currently not supported by the E-ACSL plug-in.**

## 2.3 Function contracts

*No difference with ACSL, but no terminates.*

Figure 2.6 shows grammar of function contracts. This is a simplified version of ACSL one without terminates clauses. Section 2.5 explains why E-ACSL has no terminates clause.

<i>function-contract</i>	::=	<i>requires-clause</i> *	
		<i>decreases-clause</i> <sup>?</sup>	<i>simple-clause</i> *
		<i>named-behavior</i> *	<i>completeness-clause</i> *
<i>clause-kind</i>	::=	check	admit
<i>requires-clause</i>	::=	<i>clause-kind</i> <sup>?</sup>	requires <i>pred</i> ;
<i>decreases-clause</i>	::=	decreases <i>term</i>	(for <i>id</i> ) <sup>?</sup> ;
<i>simple-clause</i>	::=	<i>assigns-clause</i>	<i>ensures-clause</i>
		<i>allocation-clause</i>	<i>abrupt-clause</i>
<i>assigns-clause</i>	::=	assigns	<i>locations</i> ;
<i>locations</i>	::=	<i>location</i>	(, <i>location</i> ) <sup>*</sup>   \nothing
<i>location</i>	::=	<i>tset</i>	
<i>ensures-clause</i>	::=	<i>clause-kind</i> <sup>?</sup>	ensures <i>pred</i> ;
<i>named-behavior</i>	::=	behavior	<i>id</i> : <i>behavior-body</i>
<i>behavior-body</i>	::=	<i>assumes-clause</i> *	<i>requires-clause</i> *
		<i>simple-clause</i> *	
<i>assumes-clause</i>	::=	assumes	<i>pred</i> ;
<i>completeness-clause</i>	::=	complete behaviors	( <i>id</i> (, <i>id</i> ) <sup>*</sup> ) <sup>?</sup> ;
		disjoint behaviors	( <i>id</i> (, <i>id</i> ) <sup>*</sup> ) <sup>?</sup> ;

Figure 2.6: Grammar of function contracts

### 2.3.1 Built-in constructs \old and \result

*No difference with ACSL.*

Figure 2.7 summarizes grammar extension of terms with \old and \result.

<i>term</i>	::=	\old ( <i>term</i> )	old value
		\result	result of a function
<i>pred</i>	::=	\old ( <i>pred</i> )	

Figure 2.7: \old and \result in terms

### 2.3.2 Simple function contracts

*No difference with ACSL.*

*\assigns is usually difficult to implement, since it requires the implementation of a memory model. Thus you would not wonder if most tools do not support it (or support it partially).*

### 2.3.3 Contracts with named behaviors

No difference with ACSL.

### 2.3.4 Memory locations and sets of terms

No difference with ACSL, but ranges and *set comprehensions* are limited in order to be finite.

Figure 2.8 describes grammar of sets of terms. The only differences with ACSL are that both lower and upper bounds of ranges are mandatory and that the predicate inside set comprehension must be guarded and bind only one variable. In that way, each set of terms is finite and their members easily identifiable.

<code>tset</code>	<code>::=</code>	<code>\emptyset</code>	empty set
		<code>tset -&gt; id</code>	
		<code>tset . id</code>	
		<code>* tset</code>	
		<code>&amp; tset</code>	
		<code>tset [ tset ]</code>	
		<code>term .. term</code>	range
		<code>\union ( tset ( , tset)* )</code>	union of locations
		<code>\inter ( tset ( , tset)* )</code>	intersection
		<code>tset + tset</code>	
		<code>( tset )</code>	
		<code>{ tset   binders ( ; pred )<sup>a</sup> }</code>	set comprehension
		<code>{ (tset ( , tset)* )<sup>b</sup> }</code>	
		<code>term</code>	implicit singleton
<code>pred</code>	<code>::=</code>	<code>\subset ( tset , tset )</code>	set inclusion
		<code>term \in tset</code>	set membership

<sup>a</sup>the given term cannot itself be a set  
<sup>b</sup>the given terms cannot themselves be a set

Figure 2.8: Grammar for sets of terms

**Example 2.5** The set `{ x | integer x; 0 <= x <= 9 || 20 <= x <= 29 }` denotes the set of all integers between 0 and 9 and between 20 and 29.

Ranges are currently only supported in memory built-ins described in Section 2.7.1, 2.13 and 2.14.

**Example 2.6** The predicate `\valid (&t[0 .. 9])` is supported and denotes that the ten first cells of the array `t` are valid. Writing the term `&t[0 .. 9]` alone, outside any memory built-in, is not yet supported.

### 2.3.5 Default contracts, multiple contracts

No difference with ACSL.

## 2.4 Statement annotations

### 2.4.1 Assertions

*No difference with ACSL.*

Figure 2.9 summarizes grammar for assertions.

<i>C-compound-statement</i>	::=	{ <i>declaration</i> * <i>statement</i> * <i>assertion</i> <sup>+</sup> }
<i>C-statement</i>	::=	<i>assertion</i> <i>statement</i>
<i>assertion-kind</i>	::=	assert   <i>clause-kind</i>
<i>assertion</i>	::=	<i>/*@ assertion-kind pred ; */</i>   <i>/*@ for id (, id)* : assertion-kind pred ; */</i>

Figure 2.9: Grammar for assertions

### 2.4.2 Loop annotations

*No difference with ACSL, but loop invariants lose their inductive nature.*

Figure 2.10 shows grammar for loop annotations. There is no syntactic difference with ACSL.

*loop assigns is usually difficult to implement, since it requires the implementation of a memory model. Thus you would not wonder if most tools do not support it (or support it partially).*

#### Loop invariants

The semantics of loop invariants is the same than the one defined in ACSL, except that they are not inductive. More precisely, if one does not take care of side effects (semantics of specifications about side effects in loop is the same in E-ACSL than the one in ACSL), a loop invariant  $I$  is valid in ACSL if and only if:

- $I$  holds before entering the loop; and
- if  $I$  is assumed true in some state where the loop condition  $c$  is also true, and if execution of the loop body in that state ends normally at the end of the body or with a "continue" statement,  $I$  is true in the resulting state.

In E-ACSL, the same loop invariant  $I$  is valid if and only if:

- $I$  holds before entering the loop; and
- if execution of the loop body in that state ends normally at the end of the body or with a "continue" statement,  $I$  is true in the resulting state.

Thus the only difference with ACSL is that E-ACSL does not assume that the invariant previously holds when one checks that it holds at the end of the loop body. In other words a loop invariant  $I$  is equivalent to put an assertion  $I$  just before entering the loop and at the very end of the loop body.

<i>statement</i>	::=	<i>/*@ loop-annot */</i> <i>while ( C-expression ) C-statement</i>   <i>/*@ loop-annot */</i> <i>for</i> <i>( C-expression ; C-expression ; C-expression )</i> <i>statement</i>   <i>/*@ loop-annot */</i> <i>do C-statement</i> <i>while ( C-expression ) ;</i>
<i>loop-annot</i>	::=	<i>loop-clause*</i> <i>loop-behavior*</i> <i>loop-variant?</i>
<i>loop-clause</i>	::=	<i>loop-invariant</i>   <i>loop-assigns</i>
<i>loop-invariant</i>	::=	<i>clause-kind?</i> <i>loop invariant pred ;</i>
<i>loop-assigns</i>	::=	<i>loop assigns locations ;</i>
<i>loop-behavior</i>	::=	<i>for id (, id)* :</i> <i>loop-clause*</i> <i>annotation for behavior id</i>
<i>loop-variant</i>	::=	<i>loop variant term ;</i>   <i>loop variant term for id ;</i> <i>variant for relation id</i>

Figure 2.10: Grammar for loop annotations

**Example 2.7** In the following, `bsearch(t,n,v)` searches for element `v` in array `t` between indices `0` and `n-1`.

```

/*@ requires n >= 0 && \valid(t+(0..n-1));
  @ assigns \nothing;
  @ ensures -1 <= \result <= n-1;
  @ behavior success:
  @ ensures \result >= 0 ==> t[\result] == v;
  @ behavior failure:
  @ assumes t_is_sorted : \forall integer k1, int k2;
  @   0 <= k1 <= k2 <= n-1 ==> t[k1] <= t[k2];
  @ ensures \result == -1 ==>
  @   \forall integer k; 0 <= k < n ==> t[k] != v;
  @*/
int bsearch(double t[], int n, double v) {
  int l = 0, u = n-1;
  /*@ loop invariant 0 <= l && u <= n-1;
    @ for failure: loop invariant
    @   \forall integer k; 0 <= k < n ==> t[k] == v ==> l <= k <= u;
    @*/
  while (l <= u) {
    int m = l + (u-l)/2; // better than (l+u)/2
    if (t[m] < v) l = m + 1;
    else if (t[m] > v) u = m - 1;
    else return m;
  }
  return -1;
}

```

In E-ACSL, this annotated function is equivalent to the following one since loop invariants are not inductive.

```

/*@ requires n >= 0 && \valid(t+(0..n-1));
  @ assigns \nothing;
  @ ensures -1 <= \result <= n-1;
  @ behavior success:
  @ ensures \result >= 0 ==> t[\result] == v;
  @ behavior failure:
  @ assumes t_is_sorted : \forall integer k1, int k2;
  @   0 <= k1 <= k2 <= n-1 ==> t[k1] <= t[k2];
  @ ensures \result == -1 ==>
  @   \forall integer k; 0 <= k < n ==> t[k] != v;
  @*/
int bsearch(double t[], int n, double v) {
  int l = 0, u = n-1;
  /*@ assert 0 <= l && u <= n-1;
    @ for failure: assert
    @   \forall integer k; 0 <= k < n ==> t[k] == v ==> l <= k <= u;
    @*/
  while (l <= u) {
    int m = l + (u-l)/2; // better than (l+u)/2
    if (t[m] < v) l = m + 1;
    else if (t[m] > v) u = m - 1;
    else return m;
  }
  /*@ assert 0 <= l && u <= n-1;
    @ for failure: assert
    @   \forall integer k; 0 <= k < n ==> t[k] == v ==> l <= k <= u;
  */
}

```



```

assertion ::= /*@ clause-kind? invariant pred ; */
           | /*@ for id (, id)* : clause-kind? invariant pred ; */

```

Figure 2.11: Grammar for general inductive invariants

```

    @*/ ;
}
return -1;
}

```

### General inductive invariant

Syntax of these kinds of invariant is shown Figure 2.11

In E-ACSL, these kinds of invariants put everywhere in a loop body is exactly equivalent to an assertion.

#### 2.4.3 Built-in construct `\at`

*No difference with ACSL, but no forward references.*

The construct `\at( $\tau$ , id)` (where `id` is a regular C label, a label added within a ghost statement or a default logic label) follows the same rule than its ACSL counterpart, except that a more restrictive scoping rule must be respected in addition to the standard ACSL scoping rule: when evaluating `\at( $\tau$ , id)` at a program point  $p$ , the program point  $p'$  denoted by `id` must be executed after  $p$  the program execution flow.

**Example 2.8** *In the following example, both assertions are accepted and valid in ACSL, but only the first one is accepted and valid in E-ACSL since evaluating the term `\at>(*p+\at(*q,Here)),L1` at L2 requires to evaluate the term `\at(*q,Here)` at L1: that is forbidden since L1 is executed before L2.*

```

/*@ requires \valid(p+(0..1));
   @ requires \valid(q);
   @*/
void f(int *p, int *q) {
    *p = 0;
    *(p+1) = 1;
    *q = 0;
    L1: *p = 2;
    *(p+1) = 3;
    *q = 1;
    L2:
    /*@ assert (\at(*p+\at(*q,L1)),Here) == 2); */
    /*@ assert (\at(*p+\at(*q,Here)),L1) == 1); */
    return ;
}

```

For the time being, `\at` can be applied to any term or predicate that uses quantified variables, let-bound variables and C variables.

**Example 2.9** The `\at` construct of the following example is supported.

```
main(void) {
  int m = 2;
  int n = 7;;
  K: ;
  n = 875;
  /*@ assert
    \let k = 3;
    \exists integer u; 9 <= u < 21 &&
    \forall integer v; -5 < v <= (u < 15 ? u + 6 : k) ==>
    \at(n + u + v > 0, K); */ ;
  return 0;
}
```

However, quantified variables that use C variables in their bounds and let-binded variables that use C variables in their definition are not yet supported.

**Example 2.10** The `\at` construct of the following example is not yet supported since the quantified variable `i` uses the C variable `n` in the definition of its upper bound.

```
/*@ ensures \forall int i; 0 <= i < n-1 ==> \old(t[i]) == t[i+1]; */
void reverse(int *t, int n) { }
```

## 2.4.4 Statement contracts

No difference with ACSL.

Figure 2.12 shows grammar of statement contracts.

<code>statement</code>	<code>::=</code>	<code>/*@ statement-contract */ statement</code>
<code>statement-contract</code>	<code>::=</code>	<code>(for id (, id)* :)? requires-clause*</code> <code>simple-clause-stmt* named-behavior-stmt*</code> <code>completeness-clause*</code>
<code>simple-clause-stmt</code>	<code>::=</code>	<code>simple-clause</code>   <code>abrupt-clause-stmt</code>
<code>named-behavior-stmt</code>	<code>::=</code>	<code>behavior id : behavior-body-stmt</code>
<code>behavior-body-stmt</code>	<code>::=</code>	<code>assumes-clause*</code> <code>requires-clause* simple-clause-stmt*</code>

Figure 2.12: Grammar for statement contracts

## 2.5 Termination

No difference with ACSL, but no terminates clauses.

### 2.5.1 Integer measures

No difference with ACSL.

### 2.5.2 General measures

*No difference with ACSL.*

### 2.5.3 Recursive function calls

*No difference with ACSL.*

### 2.5.4 Non-terminating functions

*No such feature in E-ACSL, since it is still experimental in ACSL.*

## 2.6 Logic specifications

---

*Limited to stable and computable features.*

Figure 2.13 presents grammar of logic definitions. This is the same than the one of ACSL without polymorphic definitions, lemmas, nor axiomatics.

<code>C-global-decl</code>	<code>::=</code>	<code>/*@ logic-def<sup>+</sup> */</code>
<code>logic-def</code>	<code>::=</code>	<code>logic-const-def</code>
		<code>logic-function-def</code>
		<code>logic-predicate-def</code>
<code>type-expr</code>	<code>::=</code>	<code>id</code>
<code>logic-const-def</code>	<code>::=</code>	<code>logic type-expr id = term ;</code>
<code>logic-function-def</code>	<code>::=</code>	<code>logic type-expr id parameters = term ;</code>
<code>logic-predicate-def</code>	<code>::=</code>	<code>predicate id parameters<sup>?</sup> = pred ;</code>
<code>parameters</code>	<code>::=</code>	<code>( parameter (, parameter)* )</code>
<code>parameter</code>	<code>::=</code>	<code>type-expr id</code>

Figure 2.13: Grammar for global logic definitions

### 2.6.1 Predicate and function definitions

*No difference with ACSL.*

### 2.6.2 Lemmas

*No such feature in E-ACSL: lemmas are user-given propositions. They are written usually to help theorem provers to establish validity of specifications. Thus they are mostly useful for verification activities based on deductive methods which are out of the scope of E-ACSL. Furthermore, they often requires human help to be proven, although E-ACSL targets are automatic tools.*

### 2.6.3 Inductive predicates

*No such feature in E-ACSL: inductive predicates are not computable if they really use their inductive nature.*

### 2.6.4 Axiomatic definitions

*No such feature in E-ACSL: by nature, an axiomatic is not computable.*

### 2.6.5 Polymorphic logic types

*No such feature in E-ACSL, since it is still experimental in ACSL.*

### 2.6.6 Recursive logic definitions

*No difference with ACSL.*

### 2.6.7 Higher-order logic constructions

*No such feature in E-ACSL, since it is still experimental in ACSL.*

### 2.6.8 Concrete logic types

*No such feature in E-ACSL, since it is still experimental in ACSL.*

### 2.6.9 Hybrid functions and predicates

*No difference with ACSL.*

*Hybrid functions and predicates are usually difficult to implement, since they require the implementation of a memory model (or at least to support `\at`). Thus you would not wonder if most tools do not support them (or support them partially).*

### 2.6.10 Memory footprint specification: reads clause

*No such feature in E-ACSL, since it is still experimental in ACSL.*

### 2.6.11 Specification Modules

*No difference with ACSL.*

## 2.7 Pointers and physical addressing

---

*No difference with ACSL, but separation.*

Figure 2.14 shows the additional constructs for terms and predicates which are related to memory location.

<code>term</code>	<code>::=</code>	<code>\null</code>
		<code>\base_addr one-label<sup>?</sup> ( term )</code>
		<code>\block_length one-label<sup>?</sup> ( term )</code>
		<code>\offset one-label<sup>?</sup> ( term )</code>
		<code>\allocation one-label<sup>?</sup> ( term )</code>
<code>pred</code>	<code>::=</code>	<code>\allocable one-label<sup>?</sup> ( term )</code>
		<code>\freeable one-label<sup>?</sup> ( term )</code>
		<code>\fresh two-labels<sup>?</sup> ( term, term )</code>
		<code>\valid one-label<sup>?</sup> ( location-address )</code>
		<code>\valid_read one-label<sup>?</sup> ( location-address )</code>
		<code>\separated ( location-address , location-addresses )</code>
<code>one-label</code>	<code>::=</code>	<code>{ id }</code>
<code>two-labels</code>	<code>::=</code>	<code>{ id, id }</code>
<code>location-addresses</code>	<code>::=</code>	<code>location-address ( , location-address )*</code>
<code>location-address</code>	<code>::=</code>	<code>tset</code>

Figure 2.14: Grammar extension of terms and predicates about memory

### 2.7.1 Memory blocks and pointer dereferencing

*No difference with ACSL.*

`\base_addr`, `\block_length`, `\valid`, `\valid_read` and `\offset` are usually difficult to implement, since they require the implementation of a memory model. Thus you would not wonder if most tools do not support them (or support them partially).

### 2.7.2 Separation

*No difference with ACSL.*

`\separated` is usually difficult to implement, since it requires the implementation of a memory model. Thus you would not wonder if most tools do not support it (or support it partially).

### 2.7.3 Dynamic allocation and deallocation

All these constructs are usually difficult to implement, since they require the implementation of a memory model. Thus you would not wonder if most tools do not support them (or support them partially). No difference with ACSL.

## 2.8 Sets and lists

---

### 2.8.1 Finite sets

*No difference with ACSL.*

$term$	$::=$	$[   ]$	empty list
		$[   term ( , term)^*   ]$	list of elements
		$term \hat{\ } term$	list concatenation (overloading bitwise-xor operator)
		$term *^{\ } term$	list repetition

Figure 2.15: Notations for built-in list datatype

## 2.8.2 Finite lists

*No difference with ACSL.*

Figure 2.15 shows the notations for built-in lists.

## 2.9 Abrupt termination

*No difference with ACSL.*

Figure 2.16 shows grammar of abrupt termination.

$abrupt\text{-}clause$	$::=$	$exits\text{-}clause$
$exits\text{-}clause$	$::=$	$exits\ pred ;$
$abrupt\text{-}clause\text{-}stmt$	$::=$	$breaks\text{-}clause \mid continues\text{-}clause \mid returns\text{-}clause$
$breaks\text{-}clause$	$::=$	$breaks\ pred ;$
$continues\text{-}clause$	$::=$	$continues\ pred ;$
$returns\text{-}clause$	$::=$	$returns\ pred ;$
$term$	$::=$	$\backslash exit\_status$

Figure 2.16: Grammar of contracts about abrupt terminations

## 2.10 Dependencies information

*No such feature in E-ACSL, since it is still experimental in ACSL.*

## 2.11 Data invariants

*No difference with ACSL.*

Figure 2.17 summarizes grammar for declarations of data invariants.

### 2.11.1 Semantics

*No difference with ACSL.*

<code>declaration</code>	<code>::=</code>	<code>/*@ data-inv-decl */</code>
<code>data-inv-decl</code>	<code>::=</code>	<code>data-invariant   type-invariant</code>
<code>data-invariant</code>	<code>::=</code>	<code>inv-strength<sup>?</sup> global invariant</code> <code>id : pred ;</code>
<code>type-invariant</code>	<code>::=</code>	<code>inv-strength<sup>?</sup> type invariant</code> <code>id ( C-type-name id ) = pred ;</code>
<code>inv-strength</code>	<code>::=</code>	<code>weak   strong</code>

Figure 2.17: Grammar for declarations of data invariants

### 2.11.2 Model variables and model fields

*No difference with ACSL.*

Figure 2.18 summarizes grammar for declarations of model variables and fields.

<code>declaration</code>	<code>::=</code>	<code>C-declaration</code>	
		<code>/*@ model parameter ; */</code>	model variable
		<code>/*@ model C-type-name { parameter ;<sup>?</sup> } ; */</code>	model field

Figure 2.18: Grammar for declarations of model variables and fields

## 2.12 Ghost variables and statements

---

*No difference with ACSL, but no specific construct for volatile variables.*

Figure 2.19 summarizes grammar for ghost statements which is the same than the one of ACSL.

### 2.12.1 Volatile variables

*No such feature in E-ACSL, since it is still experimental in ACSL.*

## 2.13 Initialization and undefined values

---

*No difference with ACSL.*

`\initialized` is usually difficult to implement, since it requires the implementation of a memory model. Thus you would not wonder if most tools do not support it (or support it partially).

## 2.14 Dangling pointers

---

*No difference with ACSL.*

`\dangling` is usually difficult to implement, since it requires the implementation of a memory model. Thus you would not wonder if most tools do not support it (or support it partially).

<i>C-type-qualifier</i>	::=	<i>C-type-qualifier</i>   <i>\ghost</i>	only in ghost
<i>ghost-type-specifier</i>	::=	<i>C-type-specifier</i>   <i>logic-type</i>	
<i>declaration</i>	::=	<i>C-declaration</i>   <i>/*@ ghost ghost-declaration */</i>	
<i>direct-declarator</i>	::=	<i>C-direct-declarator</i>   <i>direct-declarator</i> ( <i>C-parameter-type-list</i> <sup>?</sup> ) <i>/*@ ghost</i> ( <i>ghost-parameter-list</i> ) <i>*/</i>	ghost args
<i>postfix-expression</i>	::=	<i>C-postfix-expression</i>   <i>postfix-expression</i> ( <i>C-argument-expression-list</i> <sup>?</sup> ) <i>/*@ ghost</i> ( <i>ghost-argument-expression-list</i> ) <i>*/</i>	call with ghosts
<i>statement</i>	::=	<i>C-statement</i>   <i>statements-ghost</i>	
<i>statements-ghost</i>	::=	<i>/*@ ghost</i> <i>ghost-statement</i> <sup>+</sup> <i>*/</i>	
<i>ghost-selection-statement</i>	::=	<i>C-selection-statement</i>   <i>if ( C-expression )</i> <i>statement</i> <i>/*@ ghost else</i> <i>ghost-statement</i> <sup>+</sup> <i>*/</i>	
<i>struct-declaration</i>	::=	<i>C-struct-declaration</i>   <i>/*@ ghost</i> <i>struct-declaration</i> <i>*/</i>	ghost field

Figure 2.19: Grammar for ghost statements



## 2.15 Well-typed pointers

---

*No such feature in E-ACSL, since it is still experimental in ACSL.*

## 2.16 Logic attribute annotations

---

*No such feature in E-ACSL, since it is still experimental in ACSL.*

## 2.17 Preprocessing for ACSL

---

*No difference with ACSL.*



# Chapter 3

## Libraries

*Disclaimer:* this chapter is empty on purpose. It is left here to be consistent with the ACSL reference manual [1].



## Chapter 4

# Conclusion

This document presents an Executable ANSI/ISO C Specification Language. It provides a subset of ACSL [1] implemented [2] in the FRAMA-C platform [5] in which each construct may be evaluated at runtime. The specification language described here is intended to evolve in the future in two directions. First it is based on ACSL which is itself still evolving. Second the considered subset of ACSL may also change.





Appendix A  
Appendices

## A.1 Changes

---

### Version 1.16

- Update according to ACSL 1.16
  - **Section 2.3:** add the `check` and `admit` clause kinds.
  - **Section 2.4.1:** add the `check` and `admit` clause kinds.
  - **Section 2.4.2:** add the `check` and `admit` clause kinds.
  - **Section 2.4.2:** add the `check` and `admit` clause kinds.

### Version 1.15

- Update according to ACSL 1.15:
  - **Section 2.12:** add the `\ghost` qualifier.

### Version 1.14

- Update according to ACSL 1.14:
  - **Section 2.4.1:** add the keyword `check`.

### Version 1.13

- Update according to ACSL 1.13:
  - **Section 2.3.4:** add syntax for set membership.

### Version 1.12

- Update according to ACSL 1.12:
  - **Section 2.3.4:** add subsections for build-in lists.
  - **Section 2.4.4:** fix syntax rule for statement contracts in allowing completeness clauses.
  - **Section 2.7.1:** add syntax for defining a set by giving explicitly its element.
  - **Section 2.15:** new section.

### Version 1.9

- **Section 2.7.3:** new section.
- Update according to ACSL 1.9.

### Version 1.8

- **Section 2.3.4:** fix example 2.5.
- **Section 2.7:** add grammar of memory-related terms and predicates.



**Version 1.7**

- Update according to ACSL 1.7.
- **Section 2.7.2:** no more absent.

**Version 1.5-4**

- Fix typos.
- **Section 2.2:** fix syntax of guards in iterators.
- **Section 2.2.2:** fix definition of undefined terms and predicates.
- **Section 2.2.3:** no user-defined types.
- **Section 2.3.1:** no more implementation issue for `\old`.
- **Section 2.4.3:** more restrictive scoping rule for label references in `\at`.

**Version 1.5-3**

- Fix various typos.
- Warn about features known to be difficult to implement.
- **Section 2.2:** fix semantics of ternary operator.
- **Section 2.2:** fix semantics of cast operator.
- **Section 2.2:** improve syntax of iterator quantifications.
- **Section 2.2.2:** improve and fix example 2.3.
- **Section 2.4.2:** improve explanations about loop invariants.
- **Section 2.6.9:** add hybrid functions and predicates.

**Version 1.5-2**

- **Section 2.2:** remove laziness of operator `<==>`.
- **Section 2.2:** restrict guarded quantifications to integer.
- **Section 2.2:** add iterator quantifications.
- **Section 2.2:** extend unguarded quantifications to char.
- **Section 2.3.4:** extend syntax of set comprehensions.
- **Section 2.4.2:** simplify explanations for loop invariants and add example..

### Version 1.5-1

- Fix many typos.
- Highlight constructs with semantic changes in grammars.
- Explain why unsupported features have been removed.
- Indicate that experimental ACSL features are unsupported.
- Add operations over memory like `\valid`.
- **Section 2.2:** lazy operators `&&`, `||`, `^^`, `==>` and `<==>`.
- **Section 2.2:** allow unguarded quantification over boolean.
- **Section 2.2:** revise syntax of `\exists`.
- **Section 2.2.2:** better semantics for undefinedness.
- **Section 2.3.4:** revise syntax of set comprehensions.
- **Section 2.4.2:** add loop invariants, but they lose their inductive ACSL nature.
- **Section 2.5.2:** add general measures for termination.
- **Section 2.6.11:** add specification modules.

### Version 1.5-0

- Initial version.

## A.2 Changes in E-ACSL Implementation

---

### Version Vanadium-23.0 rc1

- **Section 2.2:** mark logic function and predicate applications as implemented.
- **Section 2.3:** mark admit clauses as implemented.

### Version Titanium-22

- **Section 2.2:** support of bitwise operations.
- **Section 2.2.7:** support of logic arrays.

### Version Scandium-21

- **Section 2.2.5:** support of rational numbers and operations.
- **Section 2.3:** remove abrupt clauses from the list of exceptions.
- **Section 2.3:** support of complete behaviors and disjoint behaviors.
- **Section 2.4.4:** remove abrupt clauses from the list of exceptions.
- **Section 2.9:** add grammar for abrupt termination.

### Version Potassium-19

- **Section 2.6:** support of logic functions and predicates.

### Version Argon-18

- **Section 2.4.3:** support of `\at` on purely logic variables.
- **Section 2.3.4:** support of ranges in memory built-ins (e.g. `\valid` or `\initialized` ).

### Version Chlorine-20180501

- **Section 2.2:** support of `\let` binding.

### Version 0.5

- **Section 2.7.3:** support of `\freeable` .

### Version 0.3

- **Section 2.4.2:** support of loop invariant.

### Version 0.2

- **Section 2.2:** support of bitwise complementation.
- **Section 2.7.1:** support of `\valid` .
- **Section 2.7.1:** support of `\block_length` .
- **Section 2.7.1:** support of `\base_addr` .
- **Section 2.7.1:** support of `\offset` .
- **Section 2.14:** support of `\initialized` .

### Version 0.1

- Initial version.



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