

Frama-C Training Session

Browsing your code dependencies

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How to better understand a C code within **Frama-C** by extracting semantic information from this code

For what purpose

- ▶ helping to start verification of an unknown code
- ▶ helping to understand results of heavier analyses
- ▶ helping heavier analyses to give better results
- ▶ helping audit activities
- ▶ helping reverse-engineering activities

In what way

- ▶ using a **battery of Frama-C plug-ins**, either syntactic or semantic



Only deduce information from a direct use of the AST

Warnings

- ▶ those here-presented use the normalised program, not the original one
- ▶ does not use advanced semantical information (for instance, the value of a variable at some statement)
- ▶ in particular, **does not handle pointers**
- ▶ some may provide incorrect results in some cases

Syntactic analyzers within Frama-C

- ▶ analysing code using program syntax only is **not the main goal of Frama-C**
- ▶ only few syntactic analyzers in Frama-C



Syntactic analyzers

what they (do not) provide

What their are good for

- ▶ getting information quickly

What their are *not* good for

- ▶ providing a big amount of useful information
- ▶ providing confidence if they may provide incorrect results



In this session, all semantic analyzers are based on abstract interpretation and the value analysis plug-in

Features

- ▶ **theoretically sound**: always provide correct results, as long as there are no soundness implementation bugs
- ▶ **handle pointers correctly**

Semantic analyzers within Frama-C

- ▶ most Frama-C plug-in are semantic analyzers



Warnings

- ▶ run the value analysis first
- ▶ may take a long time
- ▶ over-approximate the results
- ▶ all the ways to improve the efficiency/precision of the value analysis apply
- ▶ all the limitations of the value analysis also apply
- ▶ all the alarms emitted by the value analysis should be carefully examined



1. Lightweight analyzers

- ▶ Metrics
- ▶ Callgraphs
- ▶ Constant foldings
- ▶ Occurrence

*syntactic
both
both
semantic*

2. Dependencies and effects

- ▶ Functional dependencies and effects
- ▶ Imperative effects
- ▶ Operational effect
- ▶ Data scoping

*semantic
semantic
semantic
semantic*

3. Reducing code to analyse

- ▶ Slicing
- ▶ Sparecode
- ▶ Impact

*semantic
semantic
semantic*



They are either:

- ▶ syntactic analyzers; or
- ▶ semantic analyzers remaining quite precise even if the value analysis does not give so precise results

(long no
for it
C); if (0
tmp2 =
st of the

tmp2[0] = (t <= (Nb1 - 1)) ? tmp2[0] : (t <= (Nb1 - 1)) ? 0 : tmp2[0]; /* Then the second pass looks like the first one: */
tmp1[0] = 0; k = 0; k++ tmp1[0] = mc2[0][k] * tmp2[k][0]; /* The [i,j] coefficient of the matrix product MC2*TMP2, that is: *MC2[i][MC1*M1] = MC2*IM1*IMC1
i = 1; tmp1[0] >= 1; /* Final rounding: tmp2[0] is now represented on 9 bits. */ if (tmp1[0] < -256) tmp1[0] = -256; else if (tmp1[0] > 255) tmp1[0] = 255; else tmp1[0] = tmp1[0];



Give some syntactic metrics about the analyzed code.

Features

- ▶ defined and undefined functions
- ▶ number of calls to each function
- ▶ potential entry points (the never-called functions)
- ▶ number of loc
- ▶ number of conditionals, assignments, loops, calls, gotos, pointer access

Warnings

- ▶ measures are done on the normalised code, not on the original one
- ▶ does not take function pointers into account



What is it good for

- ▶ helping to measure how difficult the analyses will be
- ▶ helping to identify whether some file is missing
- ▶ helping to identify which functions have to be stubbed or specified
- ▶ helping to identify entry points of the analyzed code

How to use

- ▶ `-metrics` dumps metrics on stdout
- ▶ `-metrics-dump <f>` dumps metrics on file `f`
- ▶ also (partially) available from the GUI



Indicate the callers of each function

Features

- ▶ representation as **graphs into dot files**
- ▶ notion of **service**, a group of related functions which seems to provide common functionalities

Warning

- ▶ does not take function pointers into account

What is it good for

- ▶ helping to identify entry points of the analyzed code
- ▶ helping to discover services provided by an application
- ▶ grasping the code architecture



How to use

- ▶ `-cg <f>` dumps callgraph in dot file `f`
- ▶ `-cg-init-func <f>` adds function `f` as a root service
- ▶ from the GUI: menu **View**, then **Show Call Graph** (still experimental)

(long no
[for i <= C1; if (i
tmp2 =
st of the

tmp2[i] = (t <= (Nb1 - 1)) ? tmp2[i] : (t <= (Nb1 - 1)) ? 0 : tmp2[i] = tmp1[i]; /* Then the second pass looks like the first one: */
tmp1[0] = 0; k = 0; k++ tmp1[k] = mc2[0][k] * tmp2[k]; /* The [i,j] coefficient of the matrix product MC2*TMP2, that is: *MC2*(TMP1) = MC2*(MC1*TM1) = MC2*TM1 * MC1 =
i=1; tmp1[0] >= 1; /* Final rounding: tmp2[0] is now represented on 9 bits. */ if (tmp1[0] < -256) tmp2[0] = -256; else if (tmp1[0] > 255) tmp2[0] = 255; else tmp2[0] = tmp1[0];



Same as the syntactic callgraph...
But using the program semantics

Features

- ▶ correctly deal with function pointers

Warnings

- ▶ run the value analysis first: may take a long time

What is it good for

- ▶ computing the callgraph for codes with function pointers



How to use

- ▶ `-scg-dump` dumps the callgraph to stdout into dot format
- ▶ `-cg-init-func <f>` uses function `f` as a root service
- ▶ not available from the GUI

Warnings

- ▶ currently not the same interface as the syntactic callgraph (will be fixed soon)
- ▶ currently not exactly the same notion of service as the syntactic callgraph (will be fixed soon)



Same as the semantic callgraph...
But not represented as a graph

Feature

- ▶ display the callees of each functions

Warning

- ▶ no service computed

What is it good for

- ▶ extracting information with some external automatic tools (like grep)

How to use

- ▶ `-users` dumps the function callees on stdout



Fold all constant expressions in the code before analysis

Feature

- ▶ replace constant expressions by their results

Warning

- ▶ local propagation only: do not propagate the assignment of a constant to a left-value in the program

What is it good for

- ▶ quickly simplifying programs with lots of constant expressions
- ▶ using analysis puzzled by big constant expressions

How to use

- ▶ `-constfold` performs this analysis before all others



Propagate constant expressions in the whole program

More precisely

- ▶ generate a new program where expressions of the input program which are established as constant by the value analysis are
 - ▶ replaced by their value
 - ▶ propagated through the whole program

Features

- ▶ the output program is a compilable C code
- ▶ it has the same behaviour as the original one
- ▶ handle constant integers and pointers, even function pointers



Warning

- ▶ does not handle floating-point values yet

What is it good for

- ▶ simplifying programs with lots of constant values
- ▶ using analysis puzzled by constant expressions

How to use

- ▶ `-semantic-const-folding` propagates constants and pretty print the new source code
- ▶ `-semantic-const-fold <f1>, ..., <fn>` propagates constants only into functions `f1`, ..., `fn`
- ▶ `-cast-from-constant` replaces expressions by constants even when doing so requires a pointer cast



where variables are used

Show the uses of a variable in a program

More precisely

- ▶ highlight the left-values that may access a part of the location denoted by the selected variable

Features

- ▶ take aliasing into account
- ▶ also show uses of a C variable in logic annotations
- ▶ mainly a graphical plug-in

Warnings

- ▶ quite difficult to use in batch mode
- ▶ does not handle logic variable yet



What is it good for

- ▶ understanding a quite mysterious piece of code
- ▶ discovering some unknown aliases of the program

How to use

- ▶ **-occurrence** dumps the occurrences of each variable on stdout
- ▶ from the GUI: left panel and contextual menu

```
(long n)
{ for (i = 0; i < n; i++)
  C[i] = 0;
  tmp2 = ...
  // ...
}
```

```
tmp2[i] = (i < n) ? (i < n) : tmp2[i]; if (i < n) tmp2[i] = (i < n) ? 1 : 0; else tmp2[i] = tmp1[i]; /* Then the second pass looks like the first one:
tmp1[i] = 0; k = 0; k++ tmp1[i] = mc2[i][k] * tmp2[k]; /* The [i][k] coefficient of the matrix product MC2*TMP2, that is: *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M1 *MC1
i = 1; tmp1[i] >= 1; /* Final rounding: tmp2[i] is now represented on 9 bits. *if (tmp1[i] < -256) tmp2[i] = -256; else if (tmp1[i] > 255) tmp2[i] = 255; else tmp2[i] = tmp1[i];
```



1. Lightweight analyzers

- ▶ Metrics
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2. Dependencies and effects

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3. Reducing code to analyse

- ▶ Slicing
- ▶ Sparecode
- ▶ Impact

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Features

- ▶ several notions of input/output for functions
- ▶ several kinds of dependencies



Dependencies between inputs and outputs of functions

Definitions

- ▶ **functional output** of a function **f**: left-value that may be modified in **f** when **f** terminates
- ▶ **functional input** of a function **f**: left-value which may impact the output value of a functional output of **f**

Features

- ▶ functional outputs and inputs
- ▶ dependencies between outputs and inputs
- ▶ indicate whether the analyzer knows that an output is always modified (when the function terminates)
- ▶ ignore local variables (from the next release)



How to use

- ▶ mainly a batch plug-in
- ▶ **-deps** displays the functional dependencies for each function
- ▶ **-calldeps** displays the functional dependencies by callsite:
if a function is called several times, results are not merged

What is it good for

- ▶ providing dataflow specifications of functions
- ▶ helping to understand relations between inputs and outputs of each function
- ▶ improving precision of other analyser through **-calldeps**



- ▶ **imperative input** of a function f : left-value that may be read in f
- ▶ **imperative output** of a function f : left-value that may be written in f
- ▶ **operational input** of a function f : left-value that is read without having been previously written to, when f terminates



Features

- ▶ imperative inputs and outputs
- ▶ operational inputs

Warnings

- ▶ mainly a batch plug-in
- ▶ operational inputs are still experimental: the specification may change
- ▶ operational outputs exist but are not yet documented

(long no
for it <= 0
C1); if (0)
tmp2 =
st of the

tmp2[0] = (t <= 0 ? (N-1) : t); else if (tmp1[0] >= 0) { t <= (N-1) ? tmp2[0] : (t <= (N-1) ? 0 : tmp2[0] + tmp1[0]); } /* Then the second pass looks like the first one. */
tmp1[0] = 0; k = 0; k++ tmp1[0] = mc2[0][k] * tmp2[k][0]; /* The [i,j] coefficient of the matrix product MC2*TMP2, that is, *MC2*[t(TMP1) = MC2*[t(MC1*M1) = MC2*[M1*[MC1
i = 1; tmp1[0] >= 1; /* Final rounding: tmp2[0] is now represented on 9 bits. */ if (tmp1[0] < -256) tmp2[0] = -256; else if (tmp1[0] > 255) tmp2[0] = 255; else tmp2[0] = tmp1[0];



How to use

- ▶ **-input** displays the imperative inputs of each function; locals and function parameters are not displayed
- ▶ **-input_with_formals** same as **-input**, but displaying function parameters
- ▶ **-out** displays the imperative outputs of each function
- ▶ **-inout** displays the operational inputs of each function

```
(long n)
{ for (i = 0; i < n; i++)
  C[i] = 0;
  tmp2 = ...
  // ...
}
```

```
tmp2[0] = (i <= n-1) ? tmp2[i] : 0; if (tmp2[i] >= 0) tmp2[i] = (i <= n-1) ? tmp2[i] : 0; else tmp2[i] = tmp1[i]; /* Then the second part looks like the first one:
tmp1[0] = 0; k = 0; k++ tmp1[k] = mc2[0][k] * tmp2[k]; /* The [i][j] coefficient of the matrix product MC2*TMP1 is MC2[i][MC1*M1] = MC2[i]*M1*M1
i = 1; tmp1[0] >= 1; /* Final rounding: tmp2[0] is now represented on 9 bits: if (tmp1[0] < -256) tmp2[0] = -256; else if (tmp1[0] > 255) tmp2[0] = 255; else tmp2[0] = tmp1[0];
```



Dependencies of a given left-value $/$ at a given program point L

Features

- ▶ **show defs**: statements that may define the value of $/$ at L
- ▶ **zones**: statements that may contribute to define the value of $/$ at L
- ▶ **data scope**: statements where $/$ is guaranteed to have the same value as at L

Warning

- ▶ still experimental



What is it good for

- ▶ locally better understand what the program does
 - ▶ relations between left-values
 - ▶ where the current value of a left-value comes from
 - ▶ scope of definition of a left-value

How to use

- ▶ only available from the GUI: sub-menu **Dependencies** of the contextual menu with three entries (**Show defs**, **Zones**, **DataScope**)

```
(long n)
{ for (i = 0; i < n; i++)
  C[i] = 0;
  tmp2 = ...
  // ...
}
```

```
tmp2[i] = (i < n) ? tmp : 0; else if (tmp1[i] >= 0) { if (i < n) tmp2[i] = (i <= (n-1)) ? 0 : tmp1[i]; } // Then the second pass looks like the first one:
tmp1[0] = 0; k = 0; k++ tmp1[k] = mc2[0][k] * tmp2[k][0] / * The [i,j] coefficient of the matrix product MC2*TMP2, that is, *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M1*(MC1*
i = 1; tmp1[0][i] >= 1; /* Final rounding: tmp2[0][0] is now represented on 9 bits. */ if (tmp1[0][0] < -255) tmp2[0][0] = -255; else if (tmp1[0][0] > 255) tmp2[0][0] = 255; else tmp2[0][0] = tmp1[0][0];
}
```



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semantic
semantic
semantic



- ▶ generate a new program in a new project
- ▶ the new program is compilable
- ▶ the new program is usually shorter
- ▶ the new program is usually easier to analyze

▶ usually is not always...

Features

- ▶ generate a new program in a new project
- ▶ the new program is compilable
- ▶ the new program and the analysed one have the same behaviour according to the slicing criterion



What are the available criteria?

Criteria for code observation

- ▶ preserving effects of **statements**
- ▶ preserving the **read/write accesses** of/to left-values

Criteria for proving properties

- ▶ preserving behaviour of **assertions**
- ▶ preserving behaviour of **loop invariants**
- ▶ preserving behaviour of **loop variants**
- ▶ preserving behaviour of **threats** (emitted by the value analysis)



Pragmas

- ▶ `/*@ slice pragma ctrl; */` preserves the reachability of this control-flow point
- ▶ `/*@ slice pragma expr e; */` preserves the value of the ACSL expression `e` at this control-flow point
- ▶ `/*@ slice pragma stmt; */` preserves the effects of the next statement

How to use

- ▶ from command line options
- ▶ from the GUI: left panel and contextual menu



Each option preserves the semantics of the input program according to a specific criterion

Options of the form `-slice-criterion <f1>, ..., <fn>`

- ▶ `-slice-calls`: calls to these functions
- ▶ `-slice-return`: the return of the these functions
- ▶ `-slice-pragma`: slicing pragmas in theses functions
- ▶ `-slice-assert`: assertions of these functions
- ▶ `-slice-loop-inv`: loop invariants in these functions
- ▶ `-slice-loop-var`: loop variants in these functions
- ▶ `-slice-threat`: threats in these functions



Options of the form `-slice-criterion <v1>, ..., <vn>`

- ▶ `-slice-value` values of these left-values at the end of the entry point
- ▶ `-slice-rd` read access to these left-values
- ▶ `-slice-wr` write access to these left-values

Warning

- ▶ addresses of the left-values are evaluated at the beginning of the entry point

```
(long n)
{ for (i = 0; i < n; i++)
  C[i] = 0;
  tmp2 = ...
  // ...
}
```

```
tmp2[0] = (i < 0 ? 0 : 1); else if (tmp1[0] >= 0) { if (i < 0 ? 0 : 1) tmp2[0] = tmp1[0]; }
// ...
tmp1[0] = 0; k = 0; k++ tmp1[0] = mc2[0][k] * tmp2[k][0] / 2; // The [i][j] coefficient of the matrix product MC2*TMP2, that is, *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M1*MC1
// ...
i = 1; tmp1[0] >= 1; // Final rounding: tmp2[0][0] is now represented on 9 bits. *if (tmp1[0] < -256) tmp2[0] = -256; else if (tmp1[0] > 255) tmp2[0] = 255; else tmp2[0] = tmp1[0];
```



Custom options

- ▶ **-slicing-level <n>** specifies how to slice the callees
 - ▶ 0: never slice the called functions
 - ▶ 1: slice the callees but preserves all their functional outputs
 - ▶ 2: slice the callees but create at most 1 slice by function
 - ▶ 3: most precise slices; create as many slices as necessary

Default level is 2

- ▶ **-no-slice-undef-functions** does not slice the prototype of undefined functions (default)
- ▶ **-slice-undef-functions** slices the prototype of undefined functions
- ▶ **-slice-print** pretty prints the sliced code

Warning

- ▶ the higher the slicing level is, the slower the slicing is



- ▶ helping to extract the significant parts of a program according to your own criteria
- ▶ helping to understand where a behavior comes from
- ▶ helping analyses to give better results
- ▶ helping audit activities

Remove useless code of the program

Features

- ▶ generate a new program in a new project
- ▶ the new program is compilable
- ▶ the values assigned to the output variables of the main function are preserved in the new program
- ▶ slicing pragmas may be used to keep some statements
 - ▶ `/*@ slice pragma ctrl; */`
 - ▶ `/*@ slice pragma expr e; */`
 - ▶ `/*@ slice pragma stmt; */`



Warnings

- ▶ still experimental
- ▶ partial support of ACSL: only the annotations inside function bodies (e.g. assertions) are processed at the moment; all the others are ignored and do not appear in the new program

What is it good for

- ▶ help to discover what is useless in a program
- ▶ may improve the results of others analyzers which are puzzled by some useless code



How to use

- ▶ `-sparecode-analysis` removes statements and functions that are not useful to compute the result of the program
- ▶ `-rm-unused-globals` removes unused types and global variables
- ▶ `-sparecode-no-annot` may remove some useless code even if it changes the validity of some ACSL properties



What could be discovered
if the side effect of a statement would be revealed

More precisely

- ▶ a statement s is impacted by a statement s' iff modifying the effect of s' by another *possible* one may modify the effect of s
- ▶ an effect is *possible* iff there is an execution of the program that generates this effect. For instance, the possible effects of $z=x+y$; in $x=c?0:1$; $y=c?0:1$; $z=x+y$ are z becomes equal to 0 or 2.

Warning

- ▶ still experimental



What is it good for

- ▶ helping to understand what a statement is useful for
- ▶ helping to apprehend code changes
- ▶ helping audit activities, in particular security audits

How to use

- ▶ `-impact-pragma <f1>, ..., <fn>` computes the impact from the pragmas in functions `f1`, ..., `fn`. Only the following pragma is yet usable.

```
/*@ impact pragma stmt; */
```

- ▶ `-impact-print` dumps the result of the analysis on stdout
- ▶ from the GUI: left panel and contextual menu



Battery of Frama-C plug-ins presented

For what purpose

- ▶ helping to start verification of an unknown code
- ▶ helping to understand results of heavier analyses
- ▶ helping heavier analyses to give better results
- ▶ helping audit activities
- ▶ helping reverse-engineering activities

Browse your code dependencies more easily!

(long no
for it
C1); if (0
tmp2 =
st of the

tmp2[0] = (t <= 0) ? (t - 1) : (t + 1); if (tmp2[0] <= 0) { t = (t - 1) - 1; else tmp2[0] = tmp1[0]; } /* Then the second pass looks like the first one: "for (k = 0; k < 8; k++) tmp1[0][k] = mc2[0][k] * tmp2[k][0];" The [k][0] coefficient of the matrix product MC2*TMP2, that is, *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M1*MC1, is 1. tmp1[0][0] >= 1. */ Final rounding: tmp2[0][0] is now represented on 9 bits. *if (tmp1[0][0] < -256) m2[0][0] = -256; else if (tmp1[0][0] > 255) m2[0][0] = 255; else m2[0][0] = tmp1[0][0];

