

# Frama-C Training Session

## Browsing your code dependencies

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## How to better understand a C code within Frama-C by extracting semantic information from this code

### For what purpose

- ▶ helping to start verification of an unknown code
- ▶ helping to understand results of heavier analyses
- ▶ helping heavier analyses to give better results
- ▶ helping audit activities
- ▶ helping reverse-engineering activities

### In what way

- ▶ using a **battery of Frama-C plug-ins**, either syntactic or semantic



Only deduce information from a direct use of the AST

## Warnings

- ▶ those here-presented use the normalised program, not the original one
- ▶ does not use advanced semantical information (for instance, the value of a variable at some statement)
- ▶ in particular, **does not handle pointers**
- ▶ some may provide incorrect results in some cases

## Syntactic analyzers within Frama-C

- ▶ analysing code using program syntax only is **not the main goal of Frama-C**
- ▶ only few syntactic analyzers in Frama-C



# Syntactic analyzers

what they (do not) provide

## What their are good for

- ▶ getting information quickly

## What their are *not* good for

- ▶ providing a big amount of useful information
- ▶ providing confidence if they may provide incorrect results

long no  
for 0  
C1) if m  
tmp2  
e of the

tmp2[0] = 1 << (n-1) else if (tmp1[0] >= 1 << (n-1) - 1) tmp2[0] = (tmp1[0] >= 1 << (n-1) ? 0 : tmp1[0]) ; // Then the second part takes the first part  
tmp1[0] = 0; k = 0; k <= 5; k++) tmp1[0] += m2[0][k] \* tmp2[k]; // The [j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1 \* MC1  
l = 1; tmp1[0] >= 1; // Final rounding: tmp2[0] is now represented on 9 bits: if (tmp1[0] < -256) m2[0] = -256; else if (tmp1[0] > 255) m2[0] = 255; else tmp2[0] = tmp1[0];



In this session, all semantic analyzers are based on abstract interpretation and the value analysis plug-in

## Features

- ▶ **theoretically sound**: always provide correct results, as long as there are no soundness implementation bugs
- ▶ **handle pointers correctly**

## Semantic analyzers within Frama-C

- ▶ most Frama-C plug-in are semantic analyzers



long n...  
for 0 <= k < n  
c1[i] = m  
tmp2 =  
of the

tmp2[i] = (i <= (n-1) ? a[i] + tmp1[i]) : (i <= (n-1) ? tmp1[i] : 0); Then the second part takes for the first part  
tmp1[i] = 0; k = k + 1; tmp1[i] = m2[i][k] \* tmp2[k]; ^ The [i][j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1 \* MC1  
i = i + 1; tmp1[i] = 0; Final rounding: tmp2[i] is now represented on 3 bits: if (tmp1[i] < -256) m2[i] = -256; else if (tmp1[i] > 255) m2[i] = 255; else m2[i] = tmp1[i];

### Warnings

- ▶ run the value analysis first
- ▶ may take a long time
- ▶ over-approximate the results
- ▶ all the ways to improve the efficiency/precision of the value analysis apply
- ▶ all the limitations of the value analysis also apply
- ▶ all the alarms emitted by the value analysis should be carefully examined

```

long n;
for (i = 0; i < n; i++)
  tmp2 =
    ...

```

```

tmp2[0] = (i < (n-1)) ? tmp1[0] : 0; // The [i] coefficient of the matrix product MC2*TMP1, that is, *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M11 + MC1
tmp1[0] = 0; k = 5; k--; tmp1[0] = mc2[0][k] * tmp2[k]; // The [i] coefficient of the matrix product MC2*TMP2, that is, *MC2*(TMP2) = MC2*(MC1*M2) = MC2*M11 + MC1
i = 1; tmp1[0] >>= 1; // Final rounding, tmp2[0] is now represented on 9 bits. If (tmp1[0] < -256) tmp1[0] = -256; else if (tmp1[0] > 255) tmp1[0] = 255; else tmp1[0] =

```



## 1. Lightweight analyzers

- ▶ Metrics
- ▶ Callgraphs
- ▶ Constant foldings
- ▶ Occurrence

*syntactic*  
*both*  
*both*  
*semantic*

## 2. Dependencies and effects

- ▶ Functional dependencies and effects
- ▶ Imperative effects
- ▶ Operational effect
- ▶ Data scoping

*semantic*  
*semantic*  
*semantic*  
*semantic*

## 3. Reducing code to analyse

- ▶ Slicing
- ▶ Sparecode
- ▶ Impact

*semantic*  
*semantic*  
*semantic*



They are either:

- ▶ syntactic analyzers; or
- ▶ semantic analyzers remaining quite precise even if the value analysis does not give so precise results

long na  
for 0 <=  
ct; if (m  
tmp2 =  
of the

tmp2[0] = 1 << (N8 - 1); else if (tmp1[0]) >= 1 << (N8 - 1) tmp2[0] = 1 << (N8 - 1); else tmp2[0] = tmp1[0]; /\* Then the second part takes the first one. \*/  
tmp1[0] = 0; k = 8; k <= 16; tmp1[0] += mc2[0][k] \* tmp2[k]; /\* The [i][j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1\*MC1  
l = 1; tmp1[0] >= 1; /\* Final rounding: tmp2[0] is now represented on 9 bits. \*if (tmp1[0] < -256) tmp2[0] = -256; else if (tmp1[0] > 255) tmp2[0] = 255; else tmp2[0] = tmp1[0];



Give some syntactic metrics about the analyzed code.

### Features

- ▶ defined and undefined functions
- ▶ number of calls to each function
- ▶ potential entry points (the never-called functions)
- ▶ number of loc
- ▶ number of conditionals, assignments, loops, calls, gotos, pointer access

### Warnings

- ▶ measures are done on the normalised code, not on the original one
- ▶ does not take function pointers into account



## What is it good for

- ▶ helping to measure how difficult the analyses will be
- ▶ helping to identify whether some file is missing
- ▶ helping to identify which functions have to be stubbed or specified
- ▶ helping to identify entry points of the analyzed code

## How to use

- ▶ `-metrics` dumps metrics on stdout
- ▶ `-metrics-dump <f>` dumps metrics on file `f`
- ▶ also (partially) available from the GUI



## Indicate the callers of each function

### Features

- ▶ representation as **graphs into dot files**
- ▶ notion of **service**, a group of related functions which seems to provide common functionalities

### Warning

- ▶ does not take function pointers into account

### What is it good for

- ▶ helping to identify entry points of the analyzed code
- ▶ helping to discover services provided by an application
- ▶ grasping the code architecture



## How to use

- ▶ `-cg <f>` dumps callgraph in dot file `f`
- ▶ `-cg-init-func <f>` adds function `f` as a root service
- ▶ from the GUI: menu **View**, then **Show Call Graph** (still experimental)

long ra  
for 0 <=  
C1) if (m  
tmp2 =  
of the

tmp2[0] = 1 << (n0 - 1) else if (tmp1[0]) >> 1 << (n0 - 1) else tmp2[0] = tmp1[0]; /\* Then the second part looks like the first one. \*/  
tmp1[0] = 0; k = 5; k <= 1; tmp1[0] += m2[0][k] \* tmp2[k]; /\* The [i][j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1\*MC1  
i = 1; tmp1[0] >> 1; /\* Final rounding: tmp2[0] is now represented on 9 bits. \*/ if (tmp1[0] < -256) m2[0] = -256; else if (tmp1[0] > 255) m2[0] = 255; else tmp1[0] = tmp1[0];



Same as the syntactic callgraph...  
But using the program semantics

## Features

- ▶ correctly deal with function pointers

## Warnings

- ▶ run the value analysis first: may take a long time

## What is it good for

- ▶ computing the callgraph for codes with function pointers



## How to use

- ▶ `-scg-dump` dumps the callgraph to stdout into dot format
- ▶ `-cg-init-func <f>` uses function `f` as a root service
- ▶ not available from the GUI

## Warnings

- ▶ currently not the same interface as the syntactic callgraph (will be fixed soon)
- ▶ currently not exactly the same notion of service as the syntactic callgraph (will be fixed soon)



Same as the semantic callgraph...  
But not represented as a graph

### Feature

- ▶ display the callees of each functions

### Warning

- ▶ no service computed

### What is it good for

- ▶ extracting information with some external automatic tools (like grep)

### How to use

- ▶ `-users` dumps the function callees on stdout



Fold all constant expressions in the code before analysis

## Feature

- ▶ replace constant expressions by their results

## Warning

- ▶ local propagation only: do not propagate the assignment of a constant to a left-value in the program

## What is it good for

- ▶ quickly simplifying programs with lots of constant expressions
- ▶ using analysis puzzled by big constant expressions

## How to use

- ▶ `-constfold` performs this analysis before all others



## Propagate constant expressions in the whole program

### More precisely

- ▶ generate a new program where expressions of the input program which are established as constant by the value analysis are
  - ▶ replaced by their value
  - ▶ propagated through the whole program

### Features

- ▶ the output program is a compilable C code
- ▶ it has the same behaviour as the original one
- ▶ handle constant integers and pointers, even function pointers



## Warning

- ▶ does not handle floating-point values yet

## What is it good for

- ▶ simplifying programs with lots of constant values
- ▶ using analysis puzzled by constant expressions

## How to use

- ▶ `-semantic-const-folding` propagates constants and pretty print the new source code
- ▶ `-semantic-const-fold <f1>, ..., <fn>` propagates constants only into functions `f1`, ..., `fn`
- ▶ `-cast-from-constant` replaces expressions by constants even when doing so requires a pointer cast



## Show the uses of a variable in a program

### More precisely

- ▶ highlight the left-values that may access a part of the location denoted by the selected variable

### Features

- ▶ take aliasing into account
- ▶ also show uses of a C variable in logic annotations
- ▶ mainly a graphical plug-in

### Warnings

- ▶ quite difficult to use in batch mode
- ▶ does not handle logic variable yet



## What is it good for

- ▶ understanding a quite mysterious piece of code
- ▶ discovering some unknown aliases of the program

## How to use

- ▶ `-occurrence` dumps the occurrences of each variable on stdout
- ▶ from the GUI: left panel and contextual menu

```

long n;
for (i = 0; i < n; i++)
  tmp2 =
  ...

```

```

tmp2[i] = i * (n-1) + tmp1[i]; // i < n-1
tmp2[i] = i * (n-1) + tmp1[i]; // i < n-1
tmp1[i] = 0; k = 5; k++ tmp1[i] += mc2[i][k] * tmp2[k]; // The [i][k] coefficient of the matrix product MC2*TMP2, that is, *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M1*MC1
i = 1; tmp1[i] >= 1; // Final rounding: tmp2[i] is now represented on 9 bits: if (tmp1[i] < 256) tmp2[i] = 256; else if (tmp1[i] > 255) tmp2[i] = 255; else tmp2[i] =

```



## 1. Lightweight analyzers

- ▶ Metrics
- ▶ Callgraphs
- ▶ Constant foldings
- ▶ Occurrence

*syntactic*  
*both*  
*both*  
*semantic*

## 2. Dependencies and effects

- ▶ Functional dependencies and effects
- ▶ Imperative effects
- ▶ Operational effects
- ▶ Data scoping

*semantic*  
*semantic*  
*semantic*  
*semantic*

## 3. Reducing code to analyse

- ▶ Slicing
- ▶ Sparecode
- ▶ Impact

*semantic*  
*semantic*  
*semantic*



## Features

- ▶ several notions of input/output for functions
- ▶ several kinds of dependencies

long ra  
for 0 =>  
C1) if (m  
tmp2 =  
of the

tmp2[j] = 0; k = 0; while (tmp1[j] > 0) tmp2[j] = tmp1[j]; /\* Then the second part takes the first one. \*/  
tmp1[j] = 0; k = 0; while (tmp1[j] > 0) tmp2[j] = tmp1[j]; /\* The [j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1\*MC1  
l = 1; tmp1[j] >= 1; /\* Final rounding: tmp2[j] is now represented on 9 bits. \*M (tmp1[j]) <= 256\*m2[j] = 256; else if (tmp1[j] > 255) m2[j] = 255; else m2[j] = m2[j];



## Dependencies between inputs and outputs of functions

### Definitions

- ▶ **functional output** of a function **f**: left-value that may be modified in **f** when **f** terminates
- ▶ **functional input** of a function **f**: left-value which may impact the output value of a functional output of **f**

### Features

- ▶ functional outputs and inputs
- ▶ dependencies between outputs and inputs
- ▶ indicate whether the analyzer knows that an output is always modified (when the function terminates)
- ▶ ignore local variables (from the next release)



## How to use

- ▶ mainly a batch plug-in
- ▶ `-deps` displays the functional dependencies for each function
- ▶ `-calldeps` displays the functional dependencies by callsite:  
if a function is called several times, results are not merged

## What is it good for

- ▶ providing dataflow specifications of functions
- ▶ helping to understand relations between inputs and outputs of each function
- ▶ improving precision of other analyser through `-calldeps`



What is read, what is written,  
what is read before being written

### Definitions

- ▶ **imperative input** of a function  $f$ : left-value that may be read in  $f$
- ▶ **imperative output** of a function  $f$ : left-value that may be written in  $f$
- ▶ **operational input** of a function  $f$ : left-value that is read without having been previously written to, when  $f$  terminates



## Features

- ▶ imperative inputs and outputs
- ▶ operational inputs

## Warnings

- ▶ mainly a batch plug-in
- ▶ operational inputs are still experimental: the specification may change
- ▶ operational outputs exist but are not yet documented



### How to use

- ▶ `-input` displays the imperative inputs of each function; locals and function parameters are not displayed
- ▶ `-input_with_formals` same as `-input`, but displaying function parameters
- ▶ `-out` displays the imperative outputs of each function
- ▶ `-inout` displays the operational inputs of each function



### Dependencies of a given left-value $l$ at a given program point $L$

#### Features

- ▶ **show defs**: statements that may define the value of  $l$  at  $L$
- ▶ **zones**: statements that may contribute to define the value of  $l$  at  $L$
- ▶ **data scope**: statements where  $l$  is guaranteed to have the same value as at  $L$

#### Warning

- ▶ still experimental



## What is it good for

- ▶ locally better understand what the program does
  - ▶ relations between left-values
  - ▶ where the current value of a left-value comes from
  - ▶ scope of definition of a left-value

## How to use

- ▶ only available from the GUI: sub-menu **Dependencies** of the contextual menu with three entries (**Show defs**, **Zones**, **DataScope**)



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*semantic*  
*semantic*  
*semantic*  
*semantic*

## 3. Reducing code to analyse

- ▶ Slicing
- ▶ Sparecode
- ▶ Impact

*semantic*  
*semantic*  
*semantic*



## Features

- ▶ generate a new program in a new project
- ▶ the new program is compilable
- ▶ the new program is usually shorter
- ▶ the new program is usually easier to analyze

## Warning

- ▶ usually is not always...

long no  
for 0 <=  
ct) if (m  
tmp2 =  
of the

tmp2[j] = 0; for (k = 0; k < n; k++) tmp1[k] += m2[j][k] \* tmp2[k]; /\* The [j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP1) = MC2\*(MC1\*M1) = MC2\*M1\*MC1  
l = 1; tmp1[0] >>= 1; \*/ Final rounding: tmp2[0] is now represented on 9 bits. \*if (tmp1[0] < -256) tmp2[0] = -256; else if (tmp1[0] > 255) tmp2[0] = 255; else tmp2[0] = tmp1[0];



Specialize the program according to some user-provided criteria

### Features

- ▶ generate a new program in a new project
- ▶ the new program is compilable
- ▶ the new program and the analysed one have the same behaviour according to the slicing criterion

long n  
for (i = 0; i < n; i++)  
c[i] = 0;  
tmp2 = 0;  
for (k = 0; k < n; k++)  
tmp2 = tmp2 + c[k];

tmp2[0] = 0; for (k = 0; k < n; k++) tmp2[k] = c[k];  
The [i, j] coefficient of the matrix product MC2\*TMP2, that is, \*MC2\*(TMP2) = MC2\*(MC1\*M1) = MC2\*M1\*MC1, is 1. tmp2[0][0] >= 1. Final rounding: tmp2[0][0] is now represented on 3 bits: \*if (tmp2[0][0] < -256) tmp2[0][0] = -256; else if (tmp2[0][0] > 255) tmp2[0][0] = 255; else tmp2[0][0] = tmp2[0][0];



## What are the available criteria?

### Criteria for code observation

- ▶ preserving effects of **statements**
- ▶ preserving the **read/write accesses** of/to left-values

### Criteria for proving properties

- ▶ preserving behaviour of **assertions**
- ▶ preserving behaviour of **loop invariants**
- ▶ preserving behaviour of **loop variants**
- ▶ preserving behaviour of **threats** (emitted by the value analysis)



### Pragmas

- ▶ `/*@ slice pragma ctrl; */` preserves the reachability of this control-flow point
- ▶ `/*@ slice pragma expr e; */` preserves the value of the ACSL expression `e` at this control-flow point
- ▶ `/*@ slice pragma stmt; */` preserves the effects of the next statement

### How to use

- ▶ from command line options
- ▶ from the GUI: left panel and contextual menu



Each option preserves the semantics of the input program according to a specific criterion

Options of the form `-slice-criterion <f1>, ..., <fn>`

- ▶ `-slice-calls`: calls to these functions
- ▶ `-slice-return`: the return of the these functions
- ▶ `-slice-pragma`: slicing pragmas in theses functions
- ▶ `-slice-assert`: assertions of these functions
- ▶ `-slice-loop-inv`: loop invariants in these functions
- ▶ `-slice-loop-var`: loop variants in these functions
- ▶ `-slice-threat`: threats in these functions





## Custom options

- ▶ `-slicing-level <n>` specifies how to slice the callees
  - ▶ 0: never slice the called functions
  - ▶ 1: slice the callees but preserves all their functional outputs
  - ▶ 2: slice the callees but create at most 1 slice by function
  - ▶ 3: most precise slices; create as many slices as necessary

Default level is 2

- ▶ `-no-slice-undef-functions` does not slice the prototype of undefined functions (default)
- ▶ `-slice-undef-functions` slices the prototype of undefined functions
- ▶ `-slice-print` pretty prints the sliced code

## Warning

- ▶ the higher the slicing level is, the slower the slicing is



## What is it good for

- ▶ helping to extract the significant parts of a program according to your own criteria
- ▶ helping to understand where a behavior comes from
- ▶ helping analyses to give better results
- ▶ helping audit activities

```

long n;
for (i = 0; i < n; i++)
  tmp2[i] = 0;

```

```

tmp2[0] = 1; // (n-1) else if (tmp1[0] >= 1) // (n-1) // tmp2[0] = tmp1[0]; // Then the second pass takes the first pass
tmp1[0] = 0; k = 5; k--); tmp1[0] += mc2[0][k] * tmp2[0][k]; // The [i][j] coefficient of the matrix product MC2*TMP2, that is, *MC2*(TMP1) = MC2*(MC1*M1) = MC2*M1*MC1
l = 1; tmp1[0] >= 1; // Final spending, tmp2[0] is now represented on 9 bits. *if (tmp1[0] < -255) tmp2[0] = -255; else if (tmp1[0] > 255) tmp2[0] = 255; else tmp2[0] = tmp1[0];

```



## Remove useless code of the program

### Features

- ▶ generate a new program in a new project
- ▶ the new program is compilable
- ▶ the values assigned to the output variables of the main function are preserved in the new program
- ▶ slicing pragmas may be used to keep some statements
  - ▶ `/*@ slice pragma ctrl; */`
  - ▶ `/*@ slice pragma expr e; */`
  - ▶ `/*@ slice pragma stmt; */`



## Warnings

- ▶ still experimental
- ▶ partial support of ACSL: only the annotations inside function bodies (e.g. assertions) are processed at the moment; all the others are ignored and do not appear in the new program

## What is it good for

- ▶ help to discover what is useless in a program
- ▶ may improve the results of others analyzers which are puzzled by some useless code

long n  
for (i = 0; i < n; i++)  
c[i] = 0;  
tmp2 =  
of the

tmp2[i] = 1; // (n-1) \* 2 + 1 = 2 \* (n-1) + 1 = 2 \* n - 1 + 1 = 2 \* n  
tmp2[i] = 0; k = 5; k = tmp2[i] \* tmp2[i]; // The [i] coefficient of the matrix product MC2 \* TMP2, that is, \*MC2\*(TMP1) = MC2\*(M1 \* M1) = MC2 \* M1 \* M1  
l = 1; tmp2[i] >= 1; // Final rounding: tmp2[i] is now represented on 9 bits: if (tmp2[i] <= 255) tmp2[i] = 255; else if (tmp2[i] > 255) tmp2[i] = 255; else tmp2[i] = 255; // End of the function





## What could be discovered if the side effect of a statement would be revealed

### More precisely

- ▶ a statement  $s$  is impacted by a statement  $s'$  iff modifying the effect of  $s'$  by another *possible* one may modify the effect of  $s$
- ▶ an effect is *possible* iff there is an execution of the program that generates this effect. For instance, the possible effects of  $z=x+y$ ; in  $x=c?0:1$ ;  $y=c?0:1$ ;  $z=x+y$  are  $z$  becomes equal to 0 or 2.

### Warning

- ▶ still experimental



## What is it good for

- ▶ helping to understand what a statement is useful for
- ▶ helping to apprehend code changes
- ▶ helping audit activities, in particular security audits

## How to use

- ▶ `-impact-pragma <f1>, ..., <fn>` computes the impact from the pragmas in functions `f1`, ..., `fn`. Only the following pragma is yet usable.  
`/*@ impact pragma stmt; */`
- ▶ `-impact-print` dumps the result of the analysis on stdout
- ▶ from the GUI: left panel and contextual menu



## Battery of Frama-C plug-ins presented

### For what purpose

- ▶ helping to start verification of an unknown code
- ▶ helping to understand results of heavier analyses
- ▶ helping heavier analyses to give better results
- ▶ helping audit activities
- ▶ helping reverse-engineering activities

Browse your code dependencies more easily!

long n  
for 0 <=  
ct; if (n  
tmp2 =  
of the

tmp2[0] = 1; for (int i = 1; i < n; i++) tmp2[i] = tmp2[i-1] \* i; // First part: calculate factorial  
tmp2[0] = 0; for (int i = 1; i < n; i++) tmp2[i] = tmp2[i-1] \* i; // Second part: calculate factorial  
tmp2[0] = 0; for (int i = 1; i < n; i++) tmp2[i] = tmp2[i-1] \* i; // Third part: calculate factorial  
tmp2[0] = 0; for (int i = 1; i < n; i++) tmp2[i] = tmp2[i-1] \* i; // Fourth part: calculate factorial

